

The DragonBeam Framework: Hardware-Protected Security Modules for In-Place Intrusion Detection

Man-Ki Yoon, Mihai Christodorescu, Lui Sha, Sibir Mohan

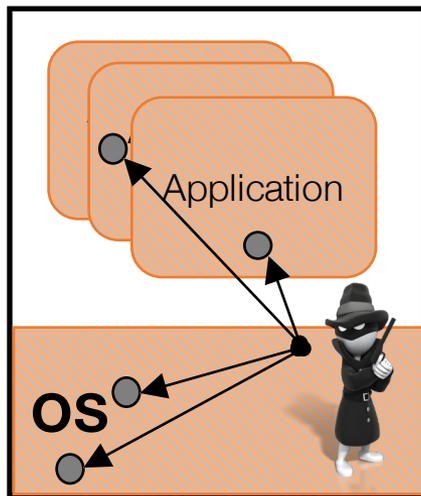
University of Illinois at Urbana-Champaign
Qualcomm Research Silicon Valley

June 6, 2016

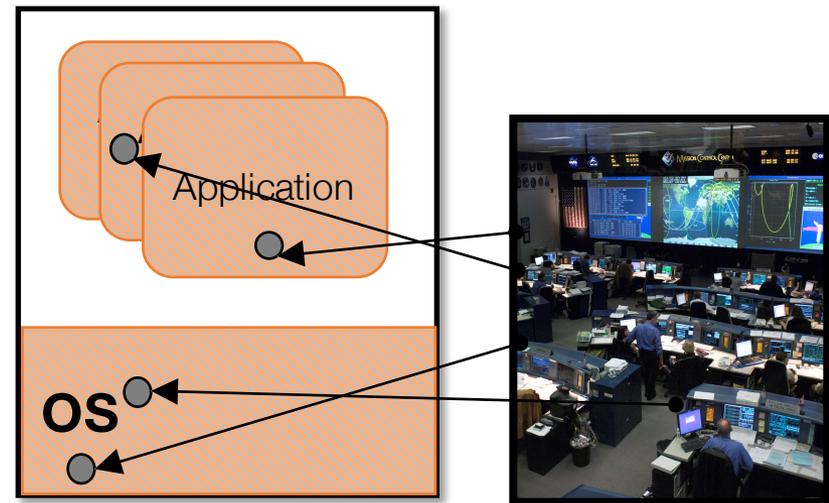


Security Monitoring

In-place Monitoring

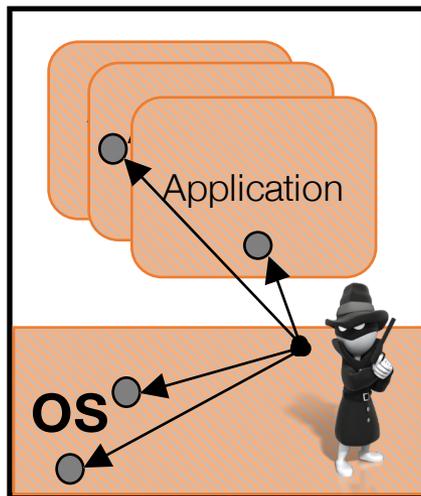


External Monitoring



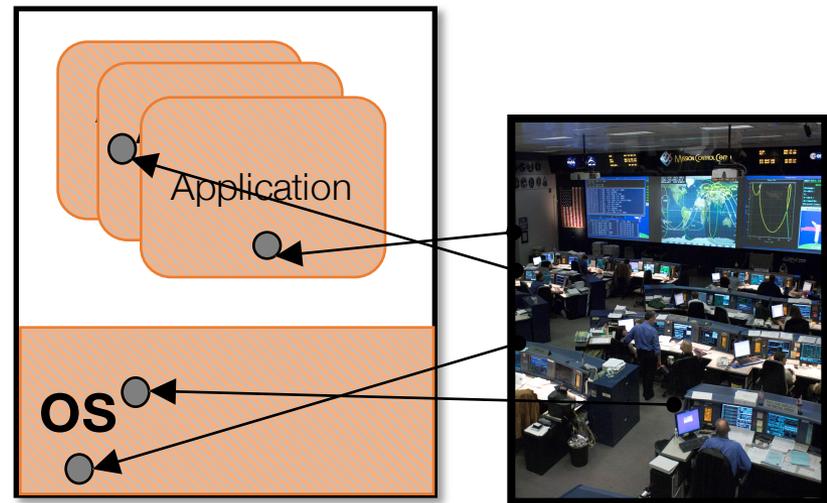
Security Monitoring

In-place Monitoring



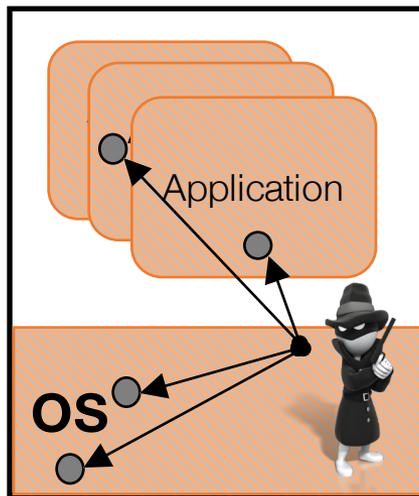
Unsafety of the monitor

External Monitoring



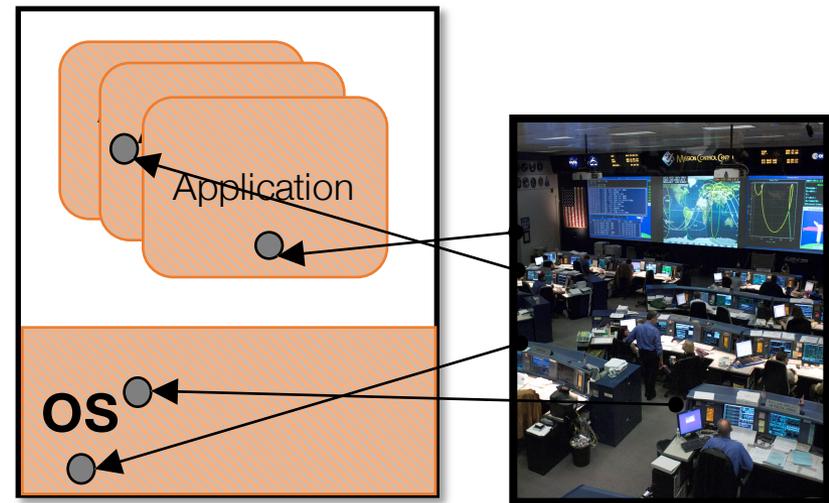
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Unsafety of the monitor

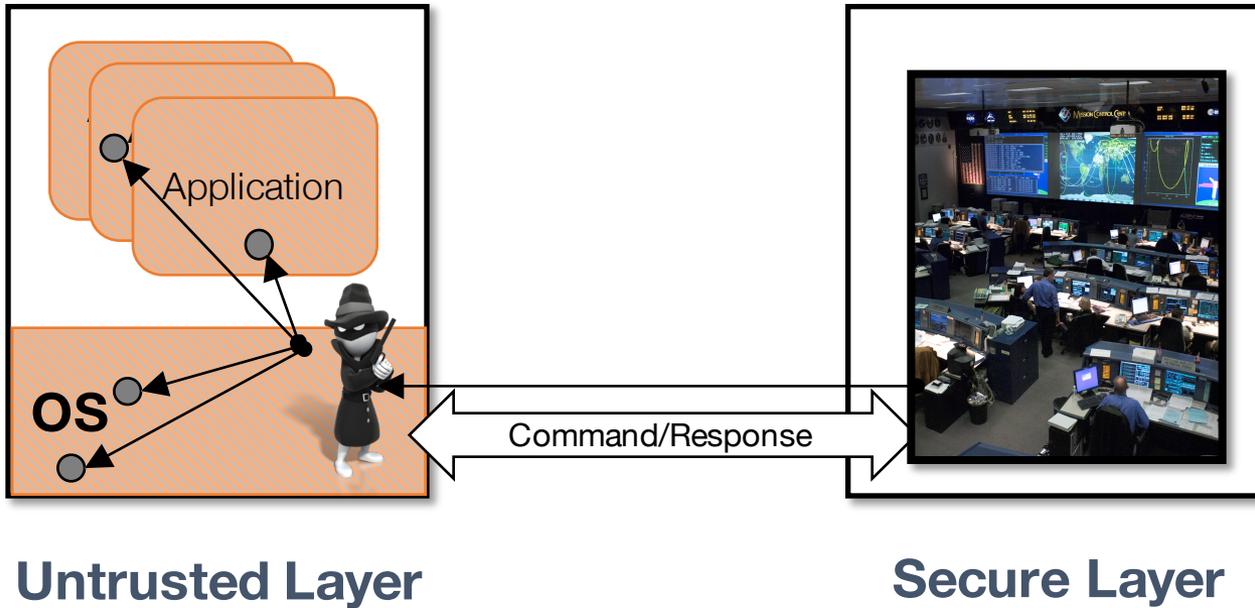
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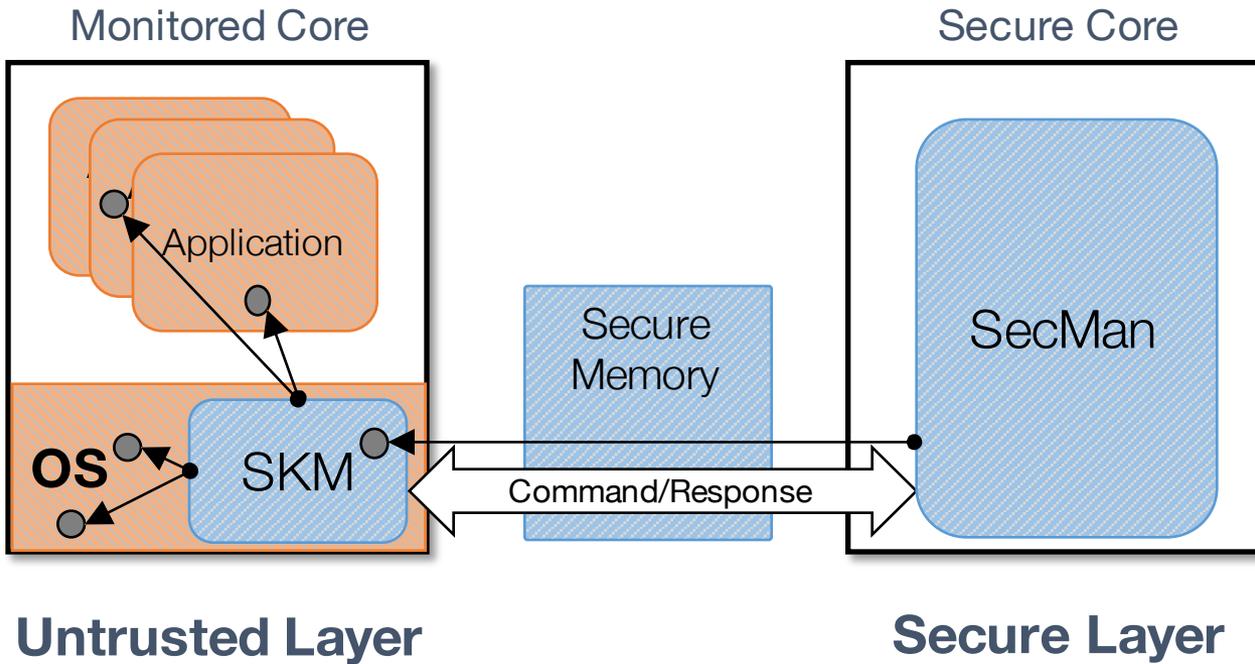
Semantic gap



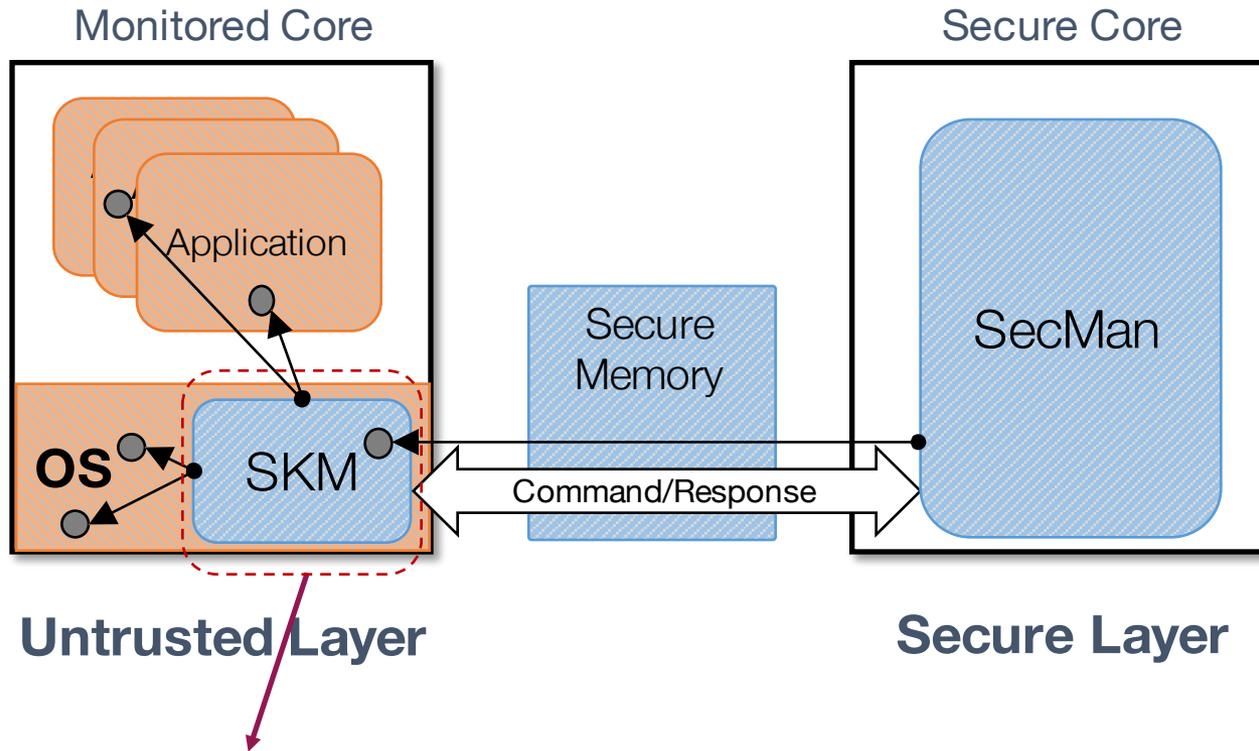
DragonBeam Framework



DragonBeam Framework



DragonBeam Framework

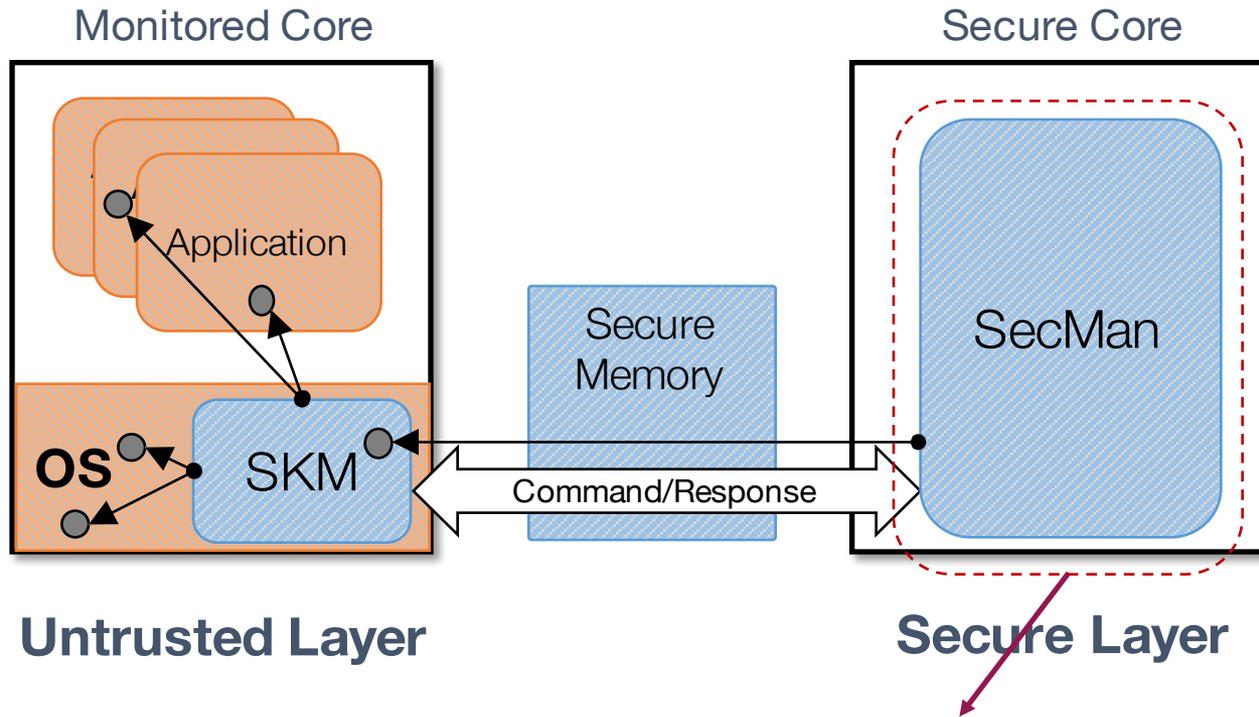


Secure Kernel Module

- Performs security monitoring operations
- Expands the observability
- Protected by Secure Core



DragonBeam Framework

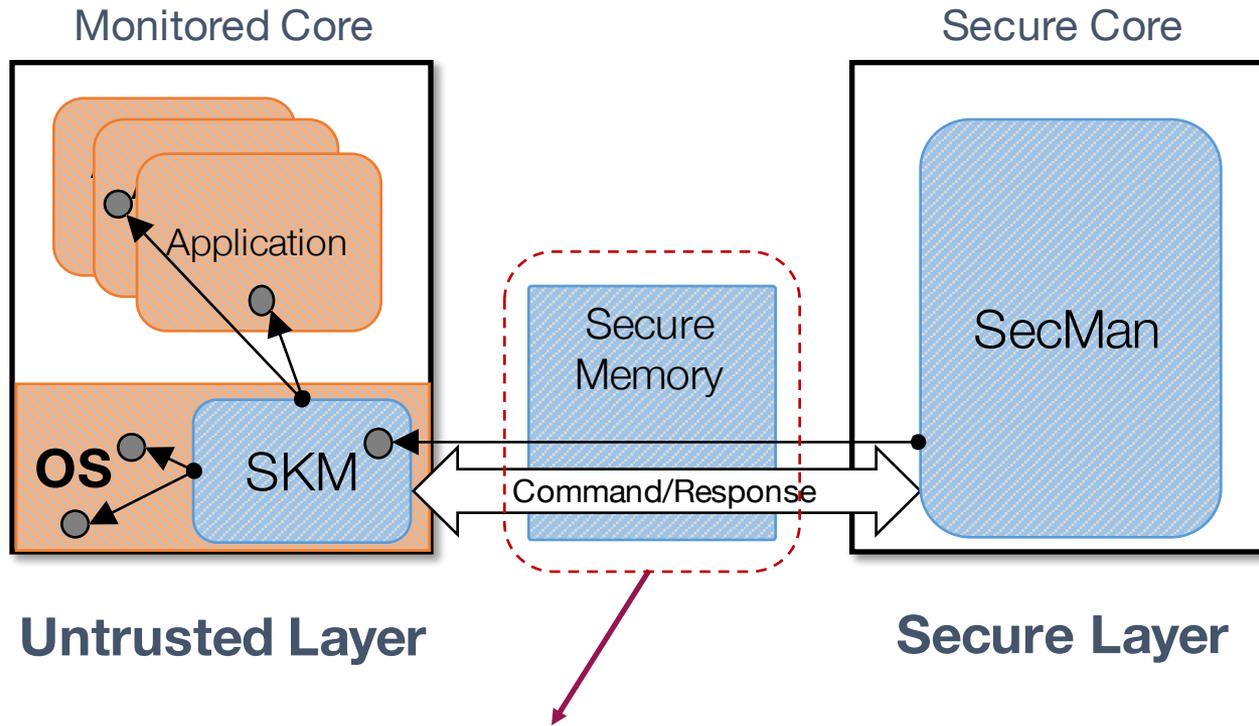


Secure Kernel Module Manager

- Commands SKM to perform security operations
- Analyzes monitored information
- Guarantees the integrity and the liveness of SKM



DragonBeam Framework

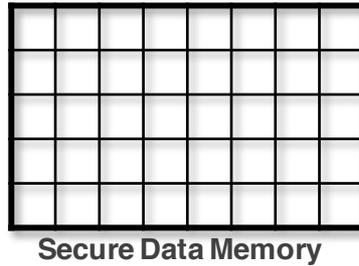
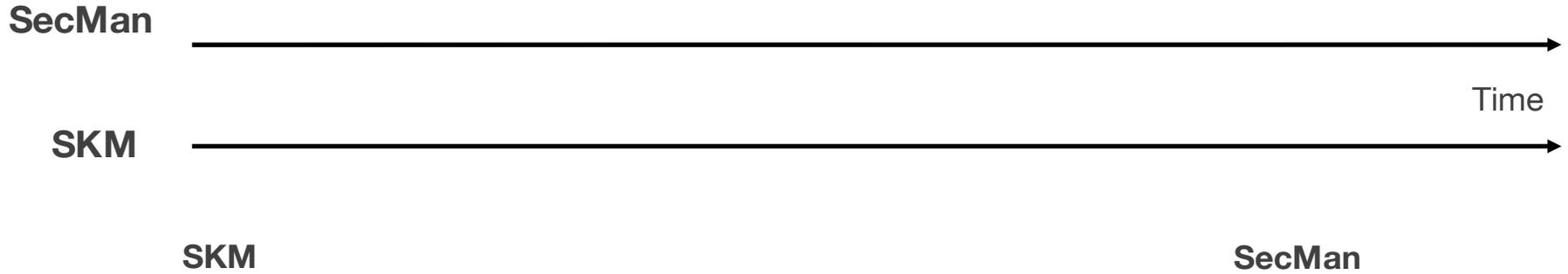


Secure Memory

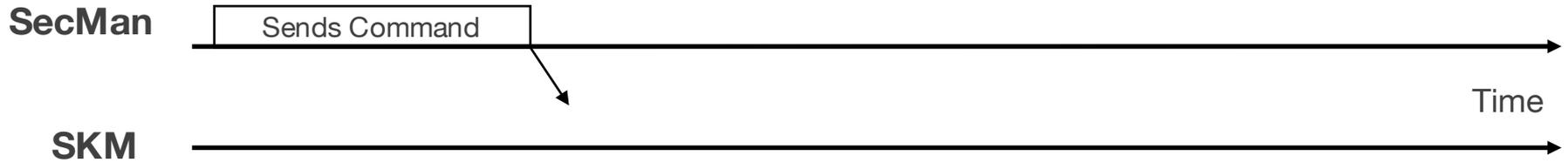
- Secure communication channel between SKM and SecMan
 - Only accessible by SKM or Secure Core
- Also hosts SecMan code/data



Example Use Case



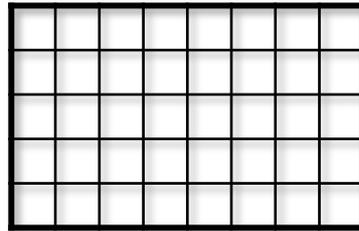
Example Use Case



SKM

SecMan

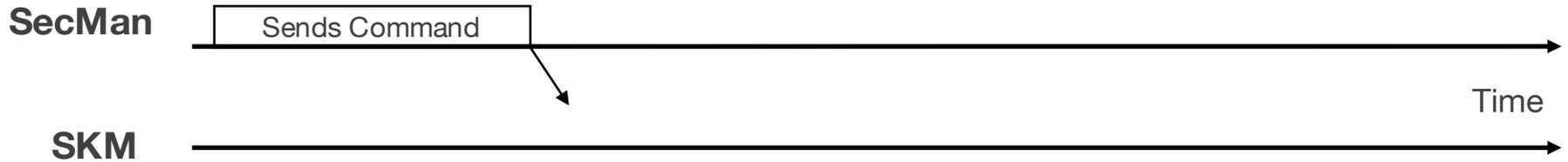
```
check_syscall_table() {  
    send_cmd(CMD_SYSCALL_TABLE);  
    settimer(TIMEOUT);  
}
```



Secure Data Memory

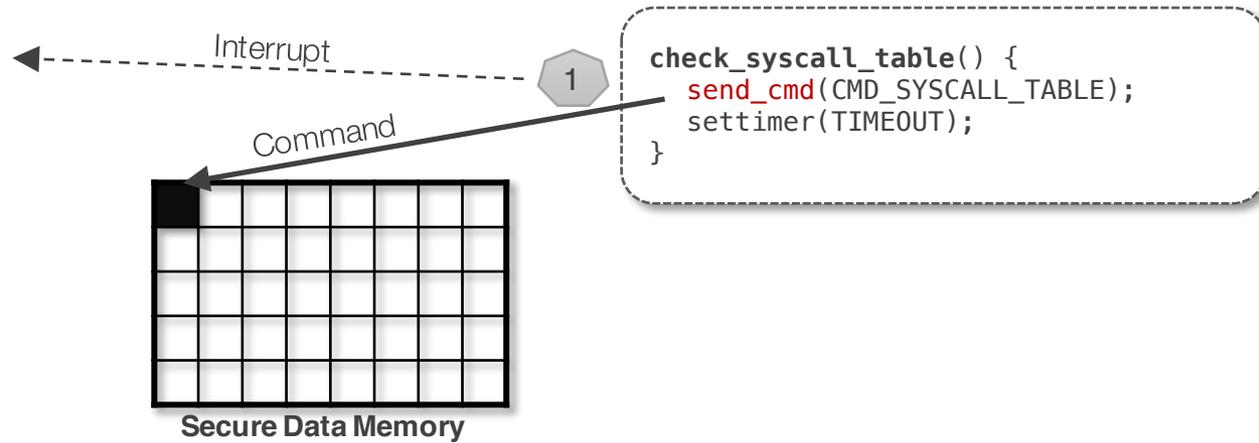


Example Use Case

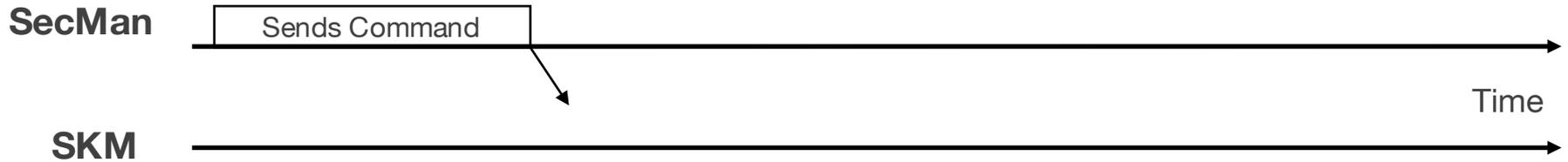


SKM

SecMan



Example Use Case



SKM

SecMan

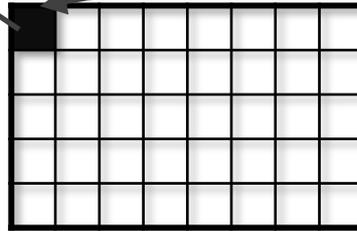
```
skm_ISR() {  
  save sp;  
  move sp to secure stack;  
  switch (*CMD) {  
    ...  
    case CMD_SYSCALL_TABLE:  
      send_syscall_table();  
      break;  
    ...  
  }  
  restore sp;  
}
```

```
check_syscall_table() {  
  send_cmd(CMD_SYSCALL_TABLE);  
  settimer(TIMEOUT);  
}
```

Interrupt

Command

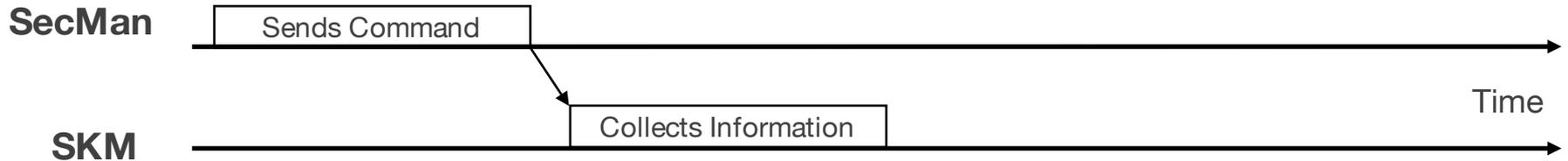
Command



Secure Data Memory



Example Use Case



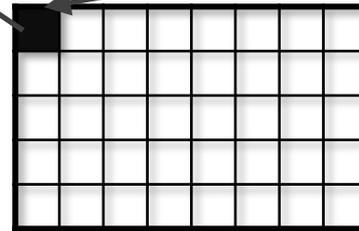
SKM

SecMan

```
skm_ISR() {  
  save sp;  
  move sp to secure stack;  
  switch (*CMD) {  
    ...  
    case CMD_SYSCALL_TABLE:  
      send_syscall_table();  
      break;  
    ...  
  }  
  restore sp;  
}
```

```
check_syscall_table() {  
  send_cmd(CMD_SYSCALL_TABLE);  
  settimer(TIMEOUT);  
}
```

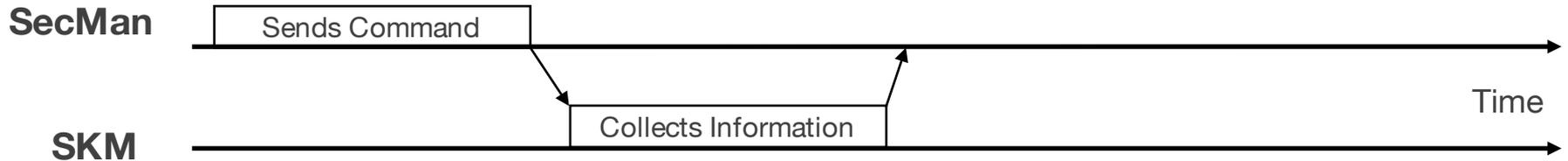
```
send_syscall_table() {  
  get cur_syscall_table;  
  for each entry i  
    write cur_syscall_table[i];  
  response_ready();  
}
```



Secure Data Memory

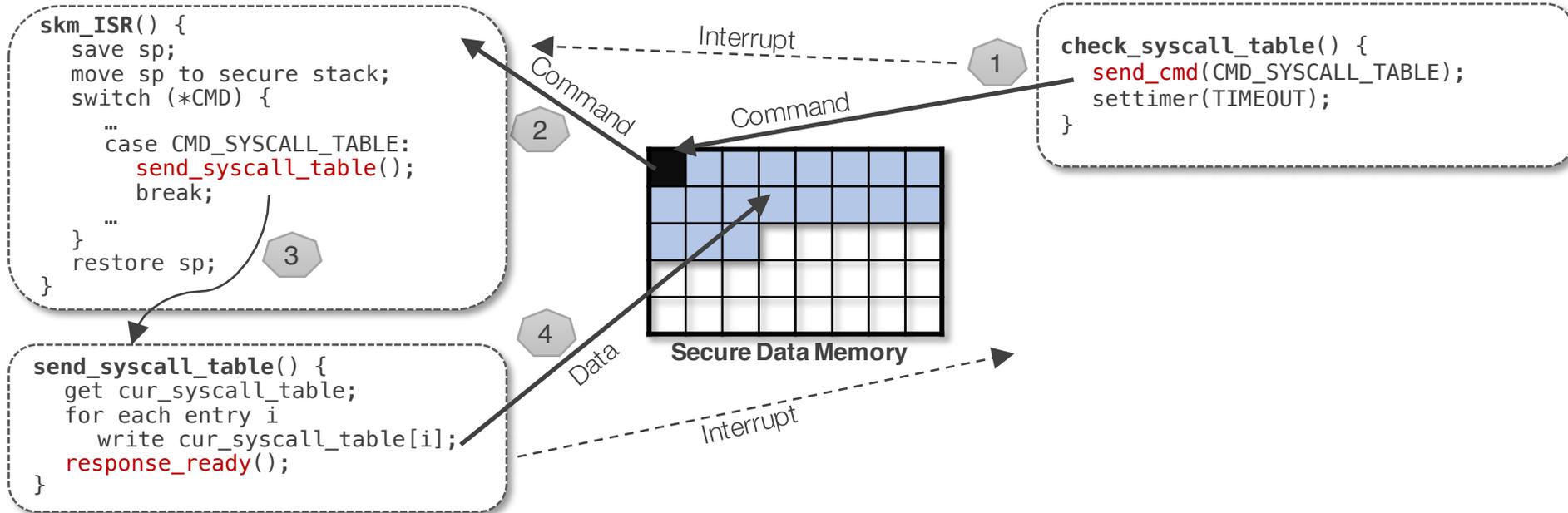


Example Use Case

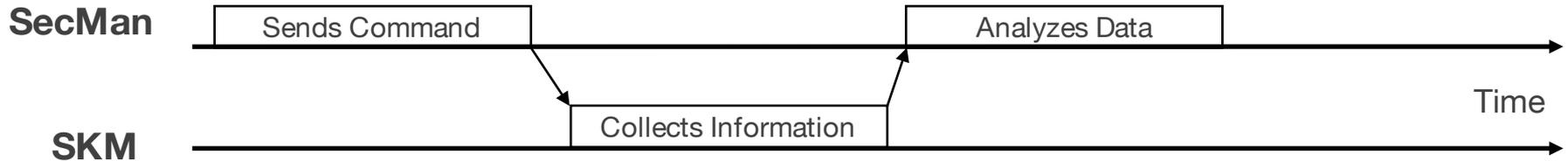


SKM

SecMan

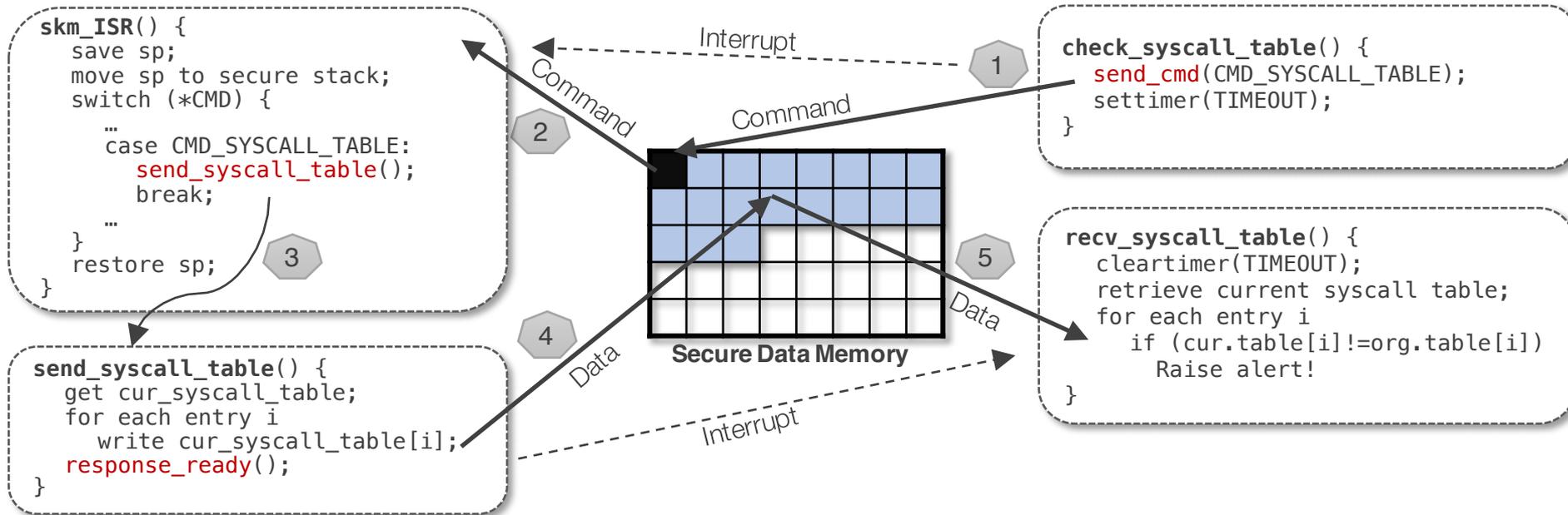


Example Use Case

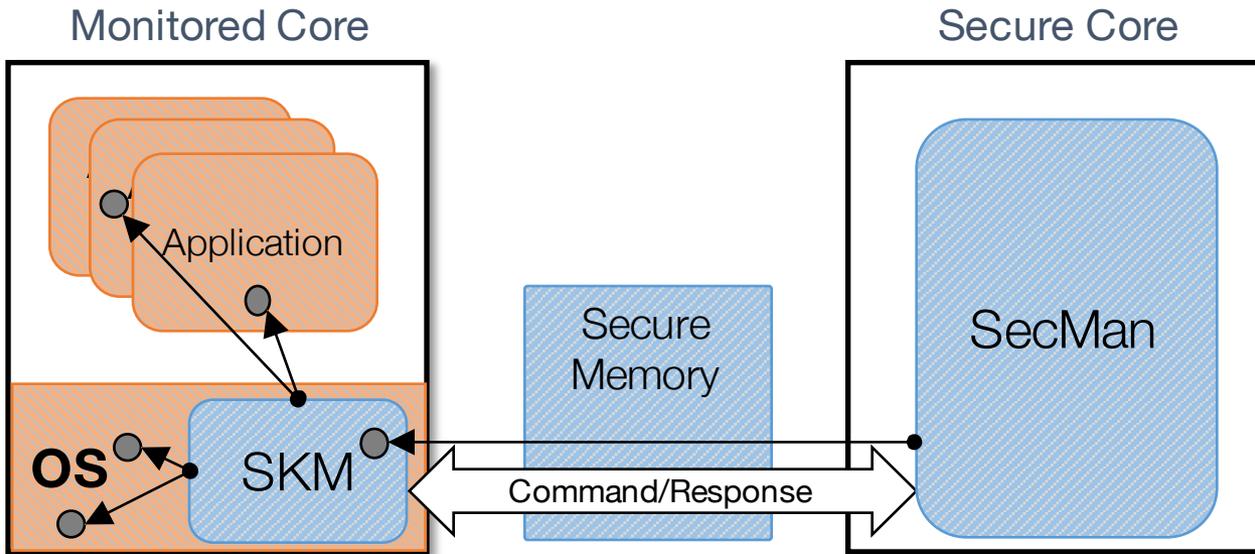


SKM

SecMan



Challenges

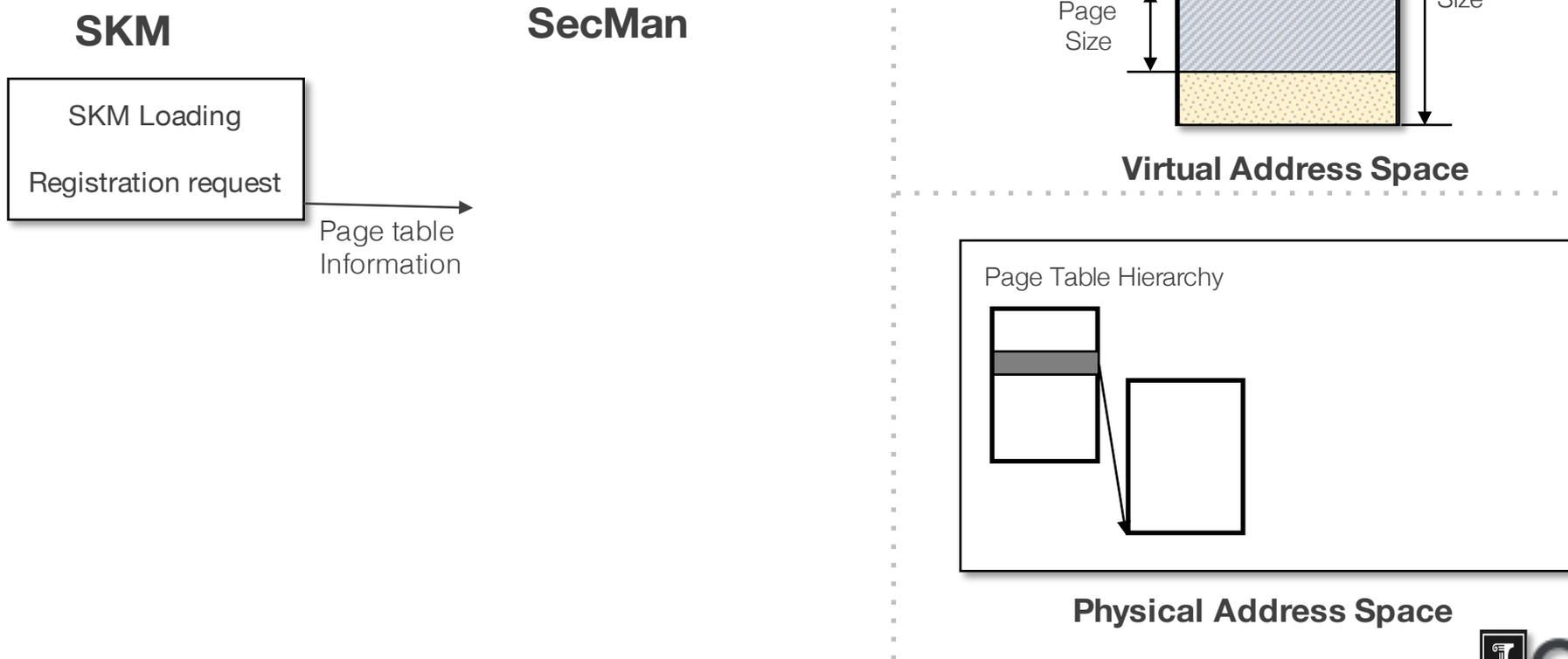


- SKM identification
- Secure memory access control
- SKM integrity and liveness guarantee

SKM Registration

- **Requested by SKM, verified by SecMan**

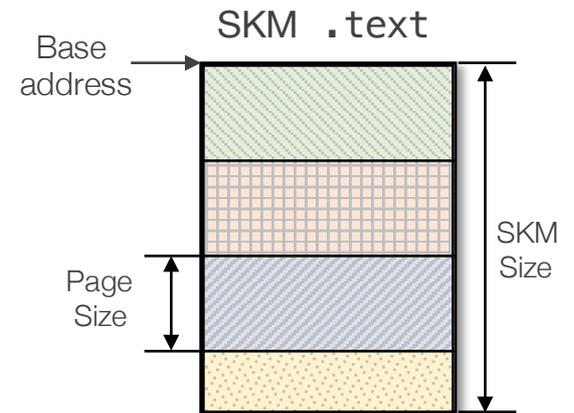
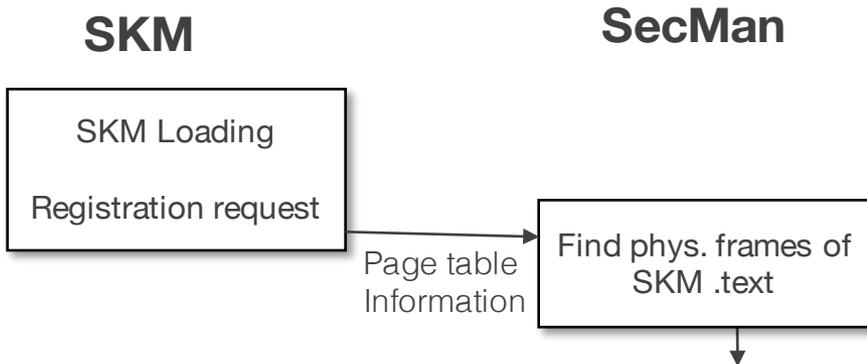
- Calculates a hash of SKM's code
- Directly from physical frames



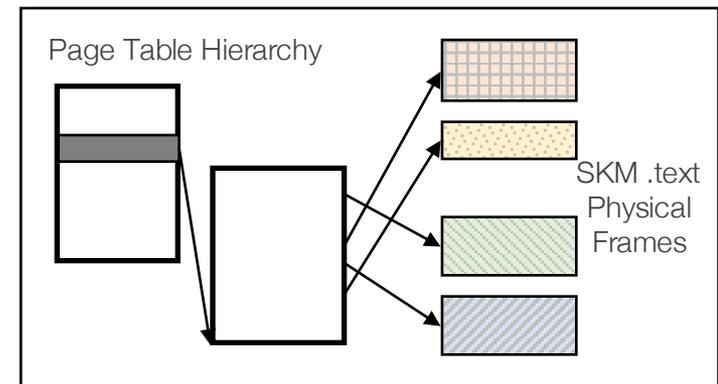
SKM Registration

- **Requested by SKM, verified by SecMan**

- Calculates a hash of SKM's code
- Directly from physical frames



Virtual Address Space

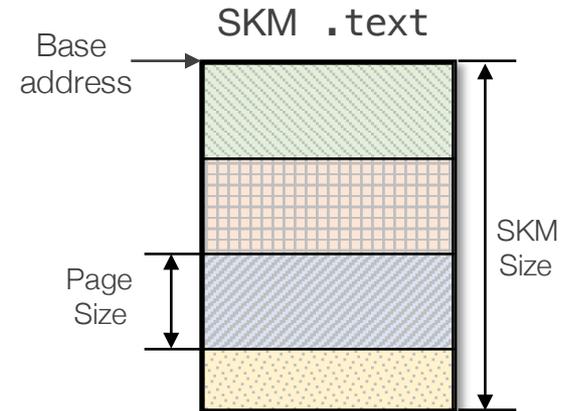
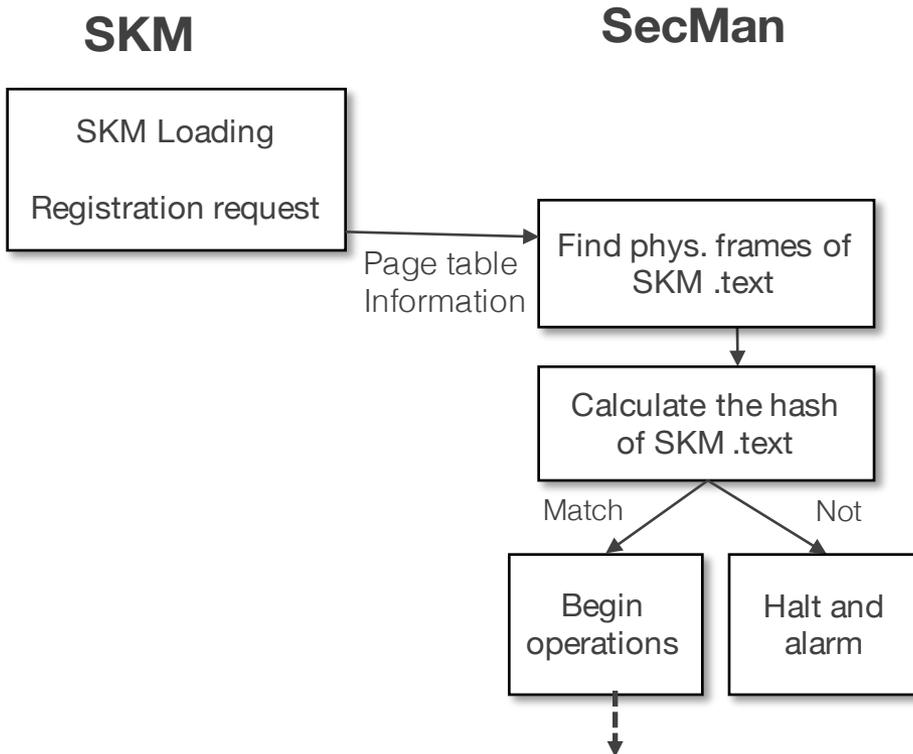


Physical Address Space

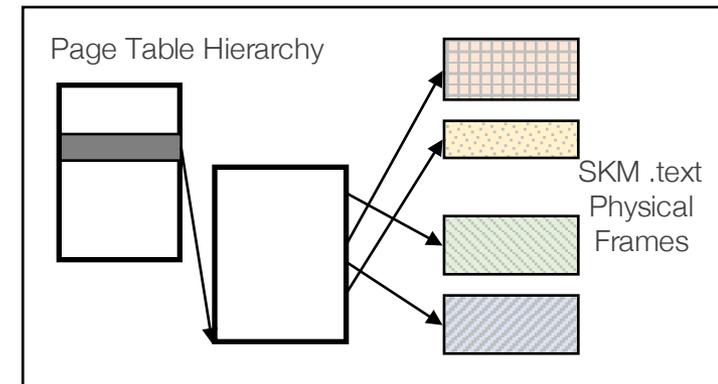


SKM Registration

- Requested by SKM, verified by SecMan
 - Calculates a hash of SKM's code
 - Directly from physical frames



Virtual Address Space



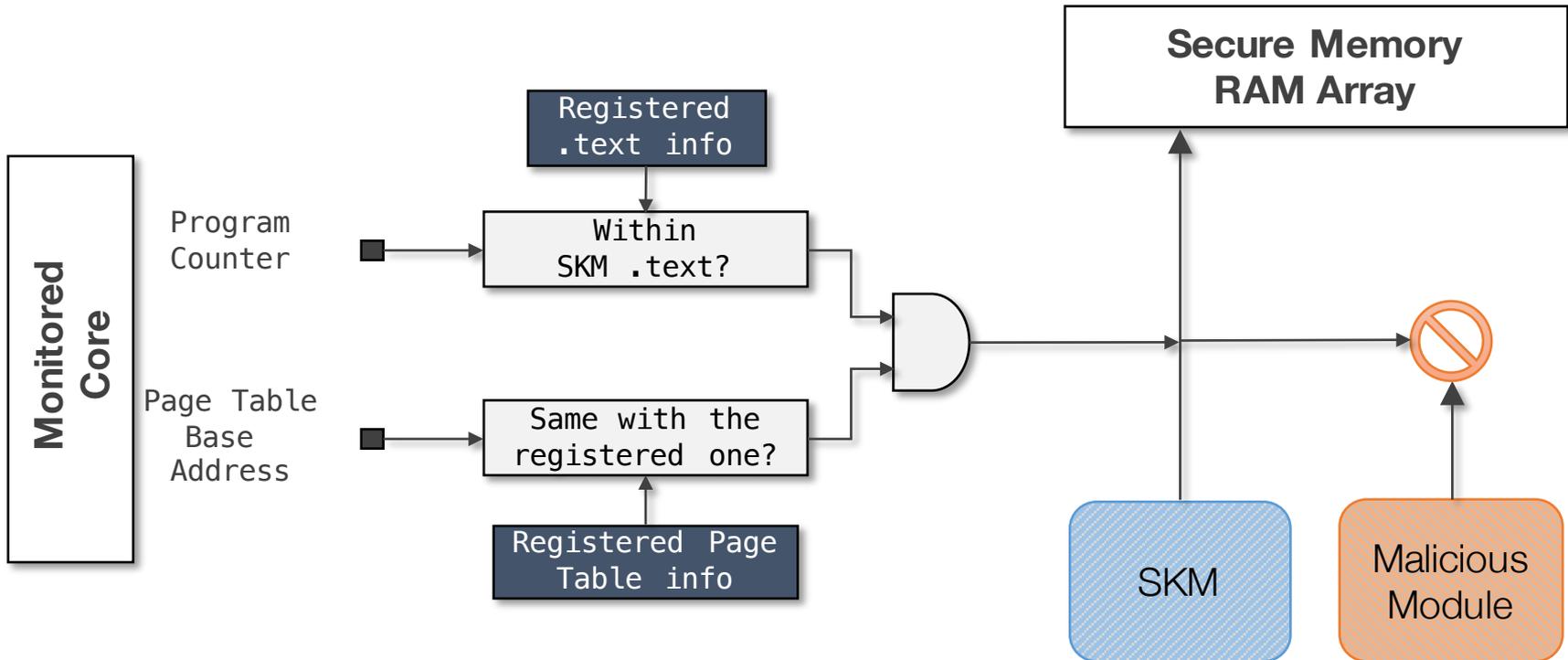
Physical Address Space



Secure Memory Access Control

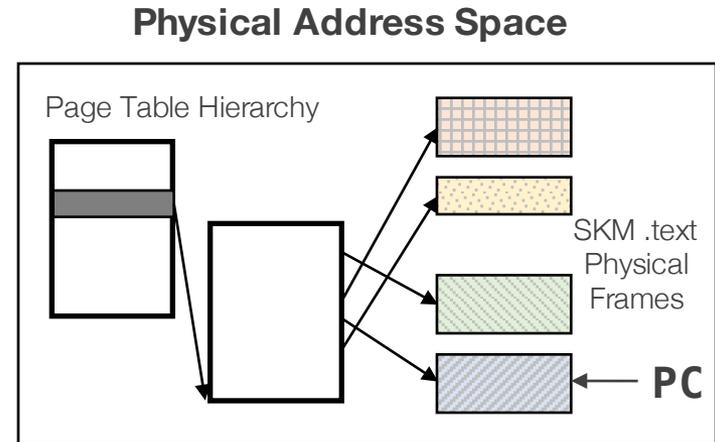
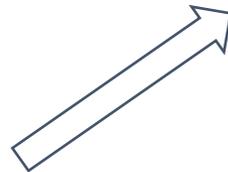
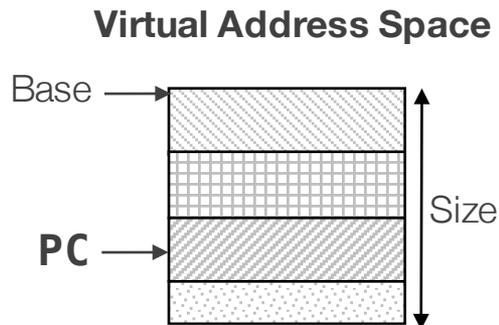
- **Who initiated memory transaction?**

- Use the current program counter and page mapping information



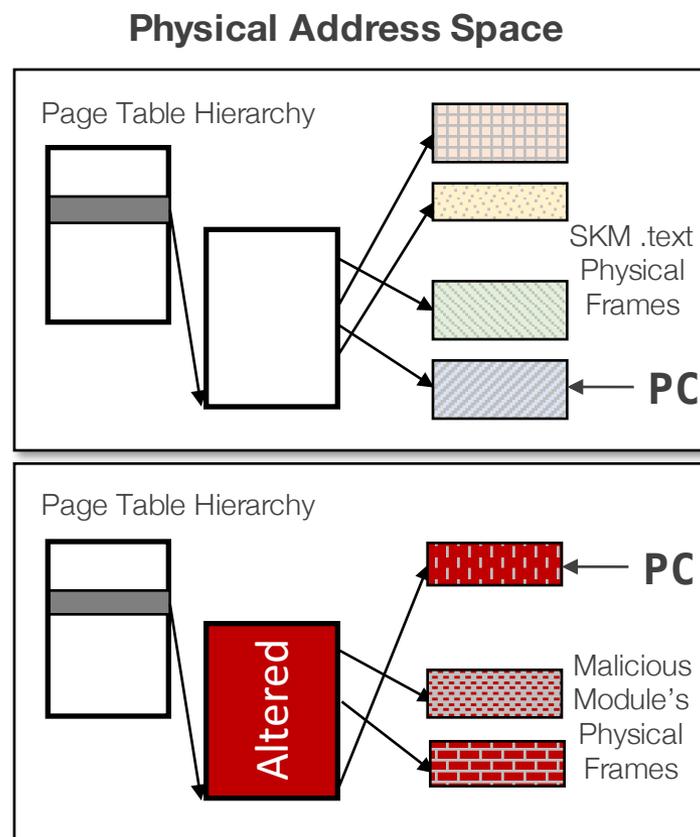
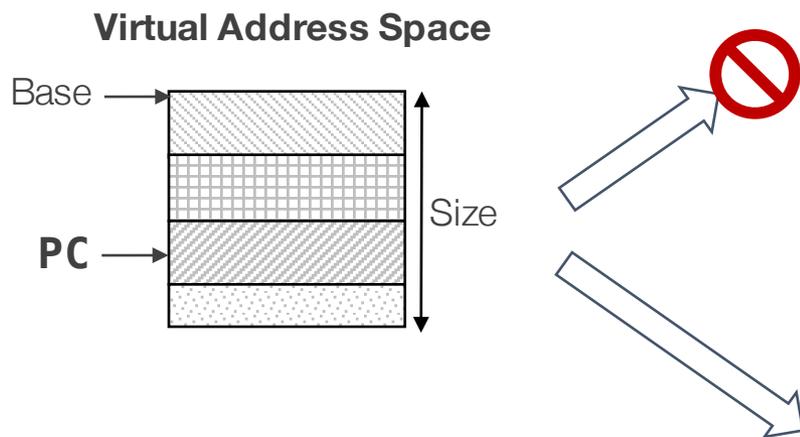
Secure Memory Access Control

- What if attacker modifies SKM's page mapping?



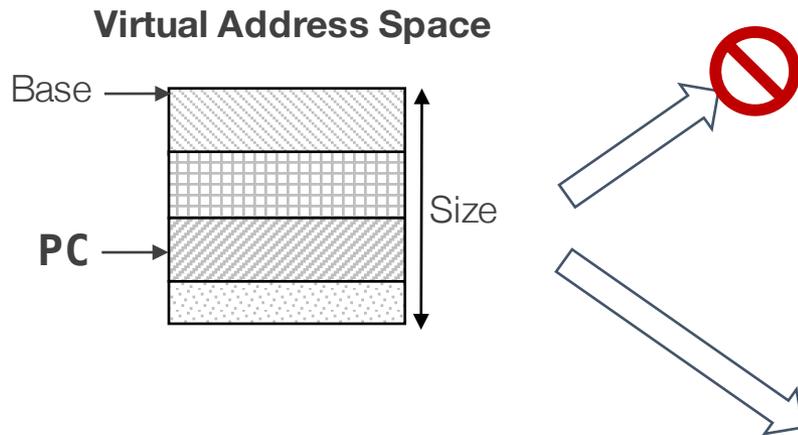
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- What if attacker modifies SKM's page mapping?

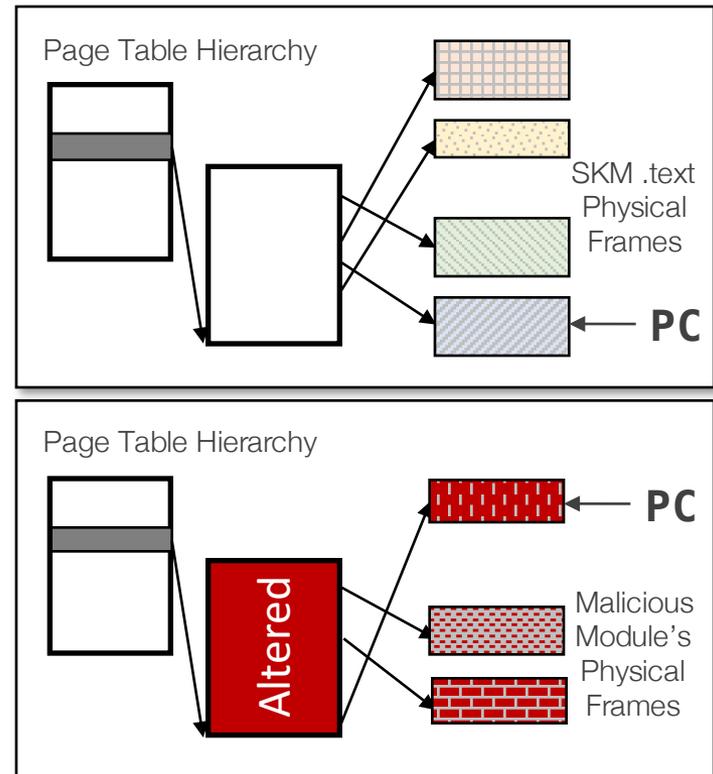


Secure Memory Access Control

- **What if attacker modifies SKM's page mapping?**
 - **Solution:** Regularly translate virt-to-phys address and verifies SKM .text hash



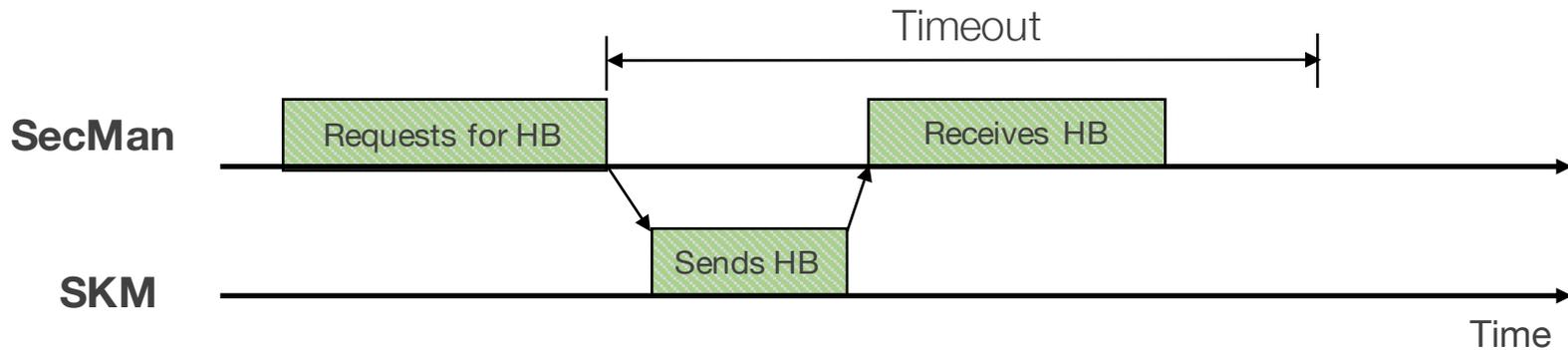
Physical Address Space



Heartbeat and Hashing

- **Heartbeat**

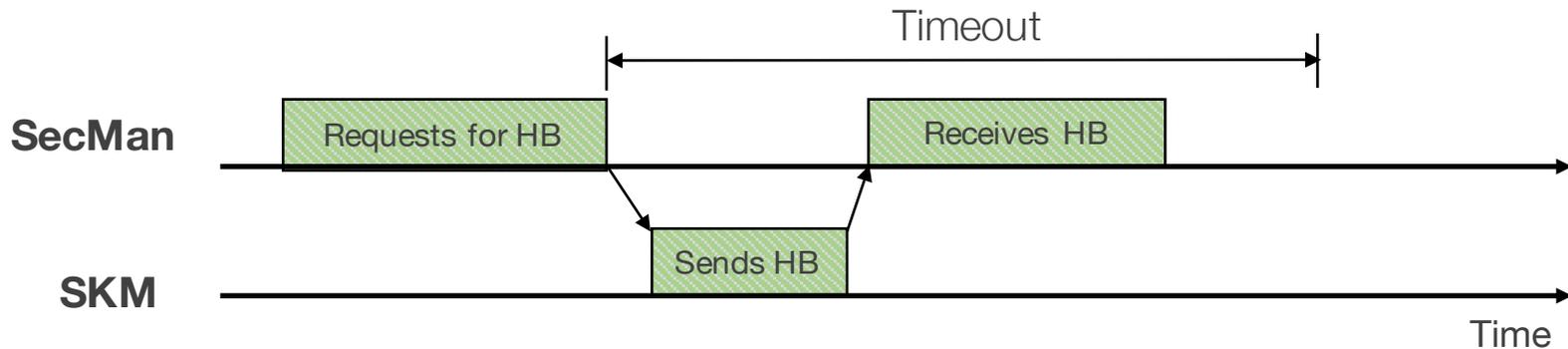
- Checks if SKM is alive
- Only SKM can respond



Heartbeat and Hashing

- **Heartbeat**

- Checks if SKM is alive
- Only SKM can respond



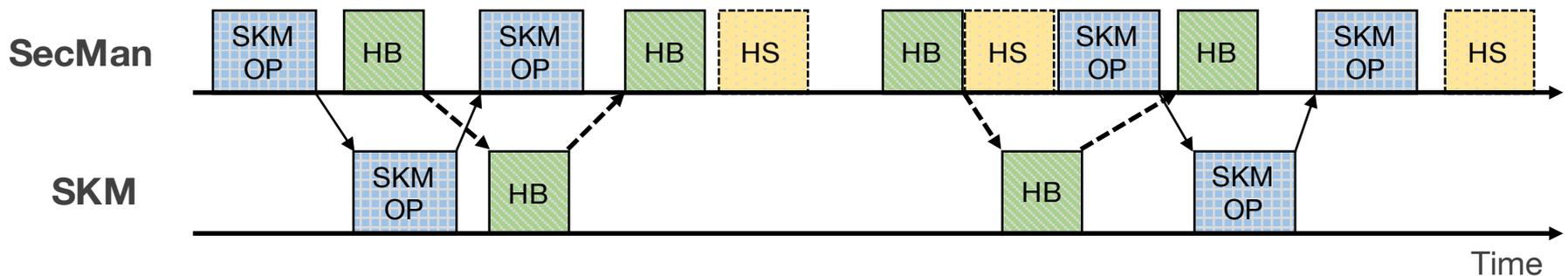
- **SKM .text hashing**

- Checks if SKM's code and page mapping have not been altered



Random Check Intervals

- To prevent TOCTTOU (Time Of Check To Time Of Use) attacks
- Attacker *cannot guess* the pattern of checks



SKM Operation
(Send/Response/Analysis)



Heartbeat
(Request/Send/Receive)



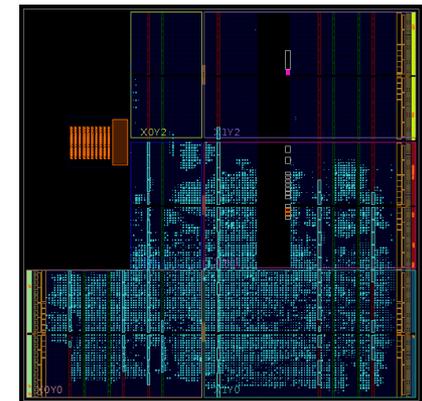
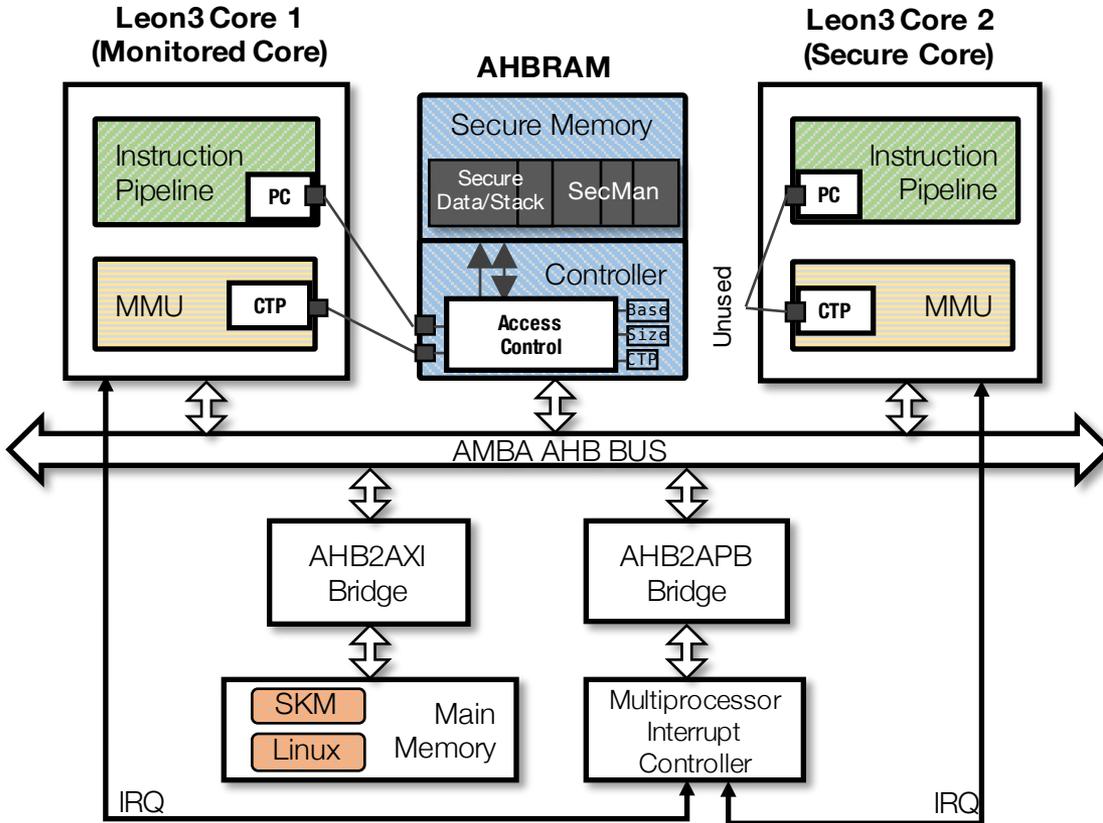
Hashing
SKM .text



Implementation

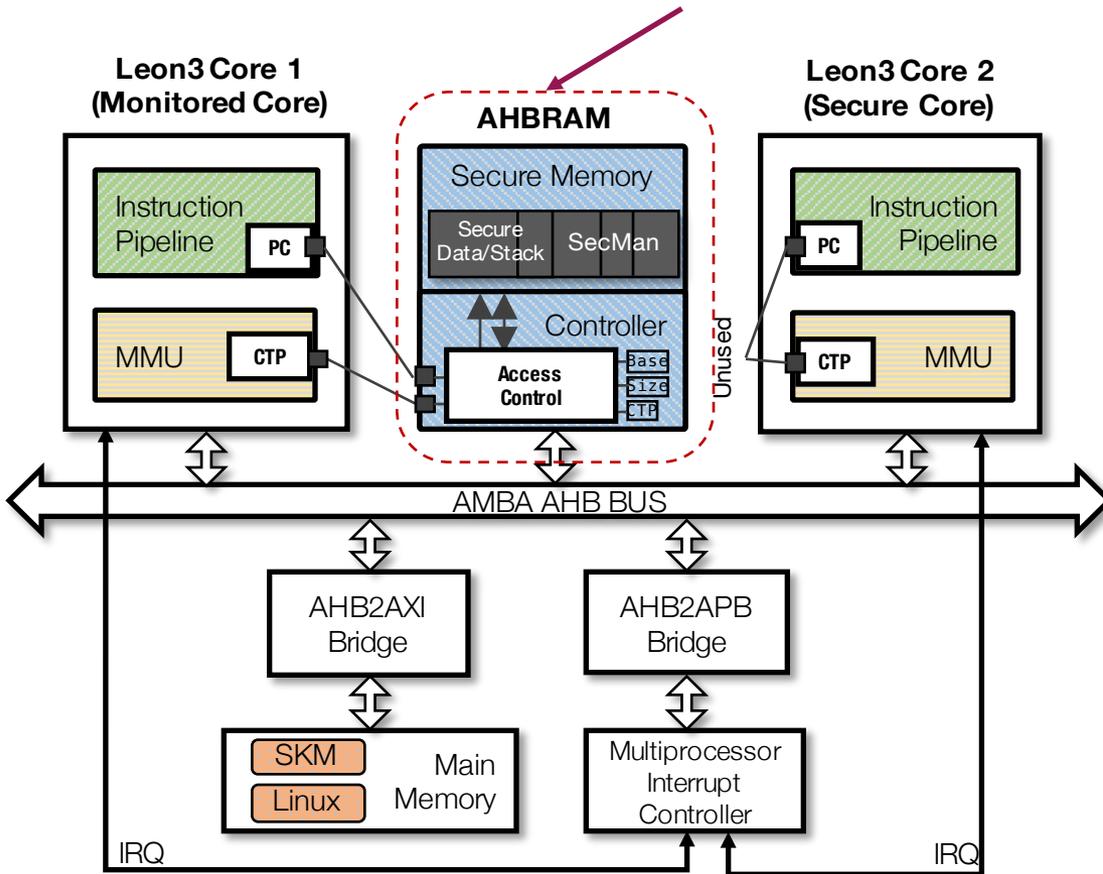
- Leon3 processor on Xilinx ZC702 FPGA

- SPARC V8, soft-core
- 83.3 MHz
- 256 MB



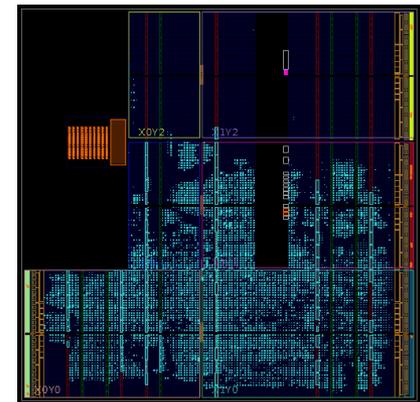
Implementation

Leon3 on-chip SRAM (128KB)



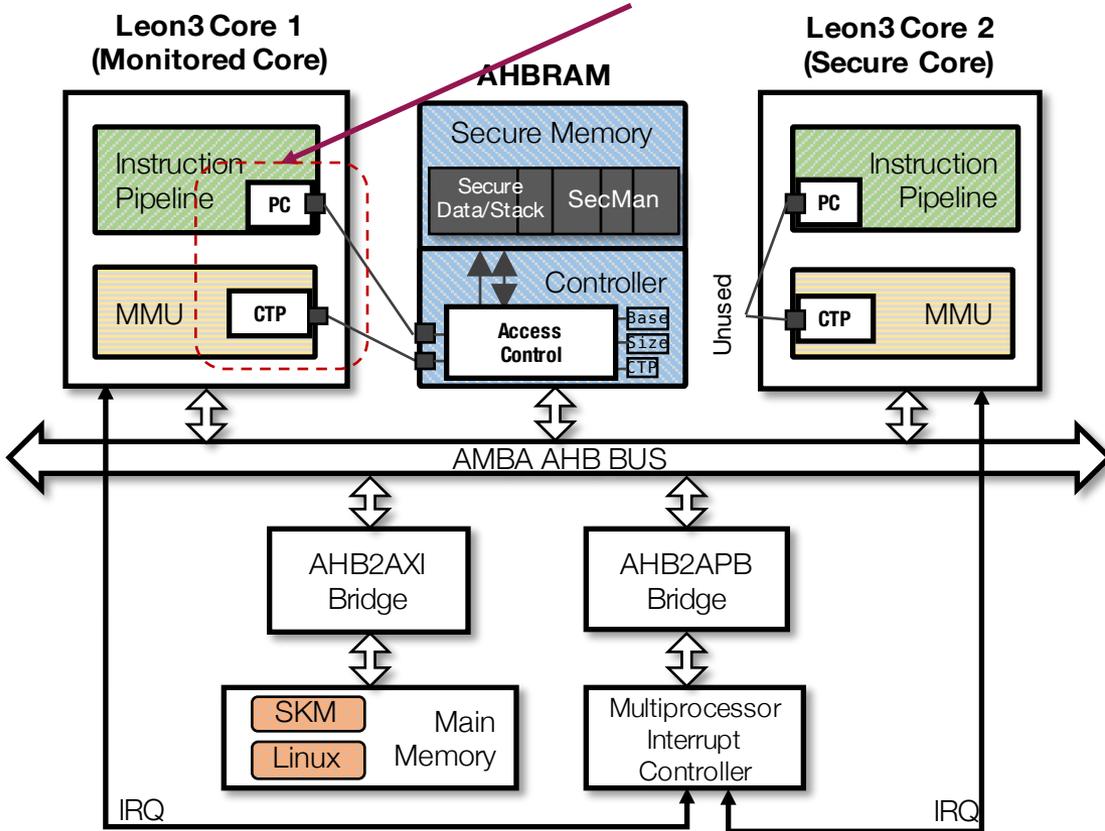
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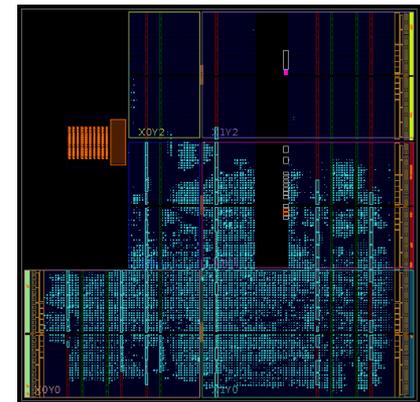
Implementation

PC and CTP for Secure Memory access control



- Leon3 processor on Xilinx ZC702 FPGA

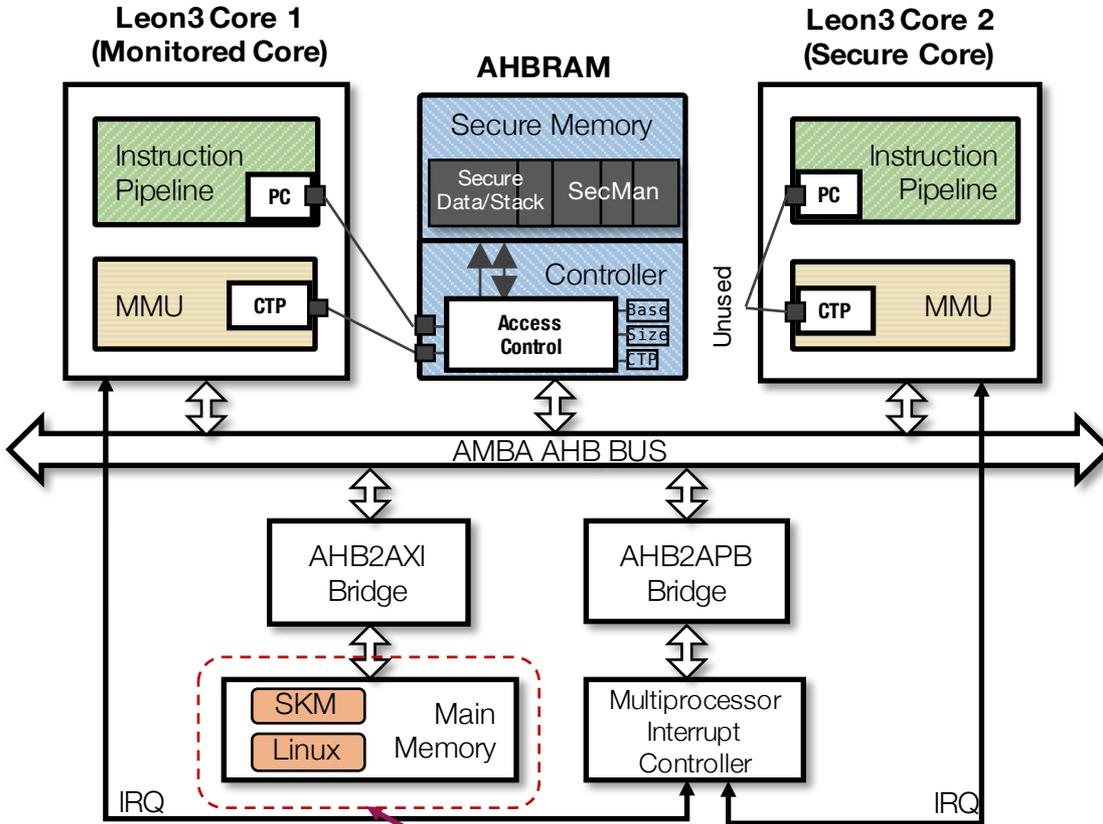
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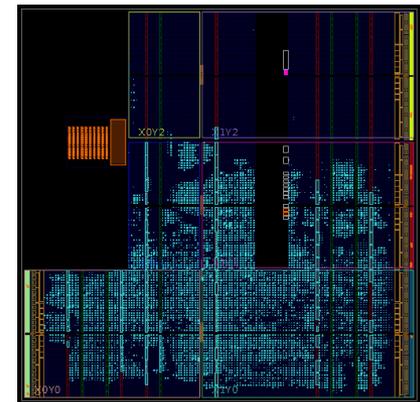
Implementation

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Unmodified Linux 3.8



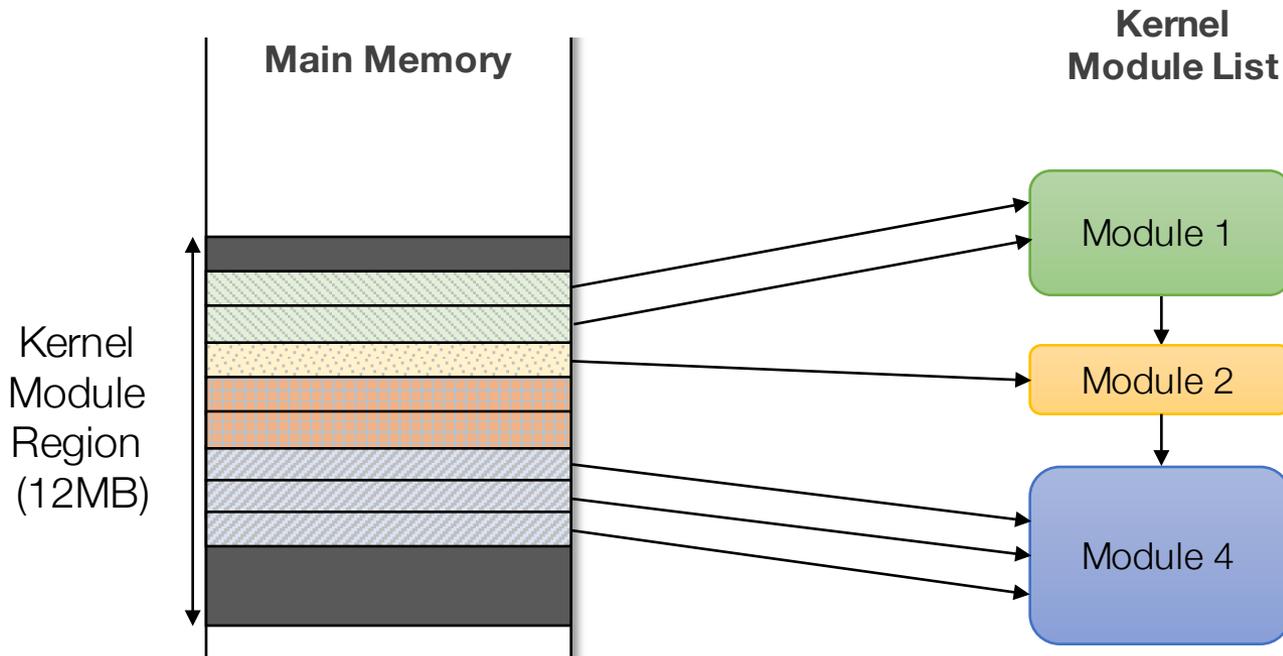
Evaluation – Use Cases

- 1) System call table integrity check
- 2) Hidden module detection



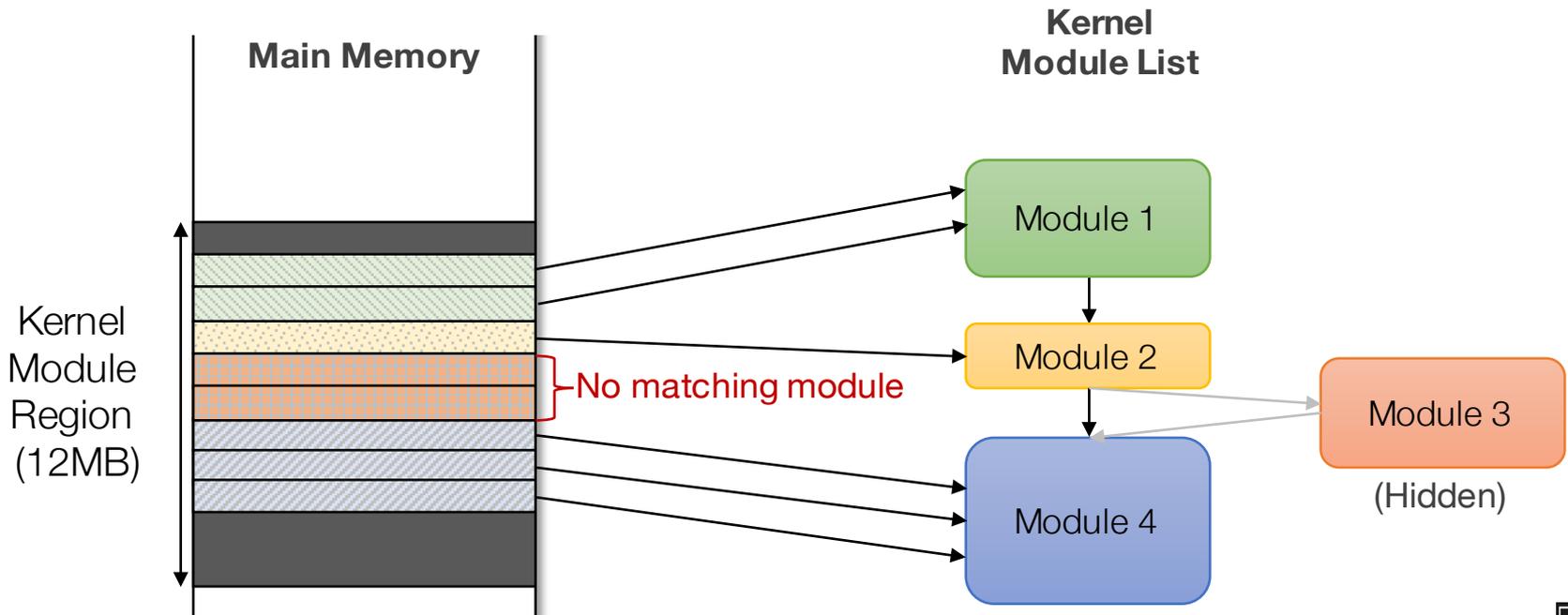
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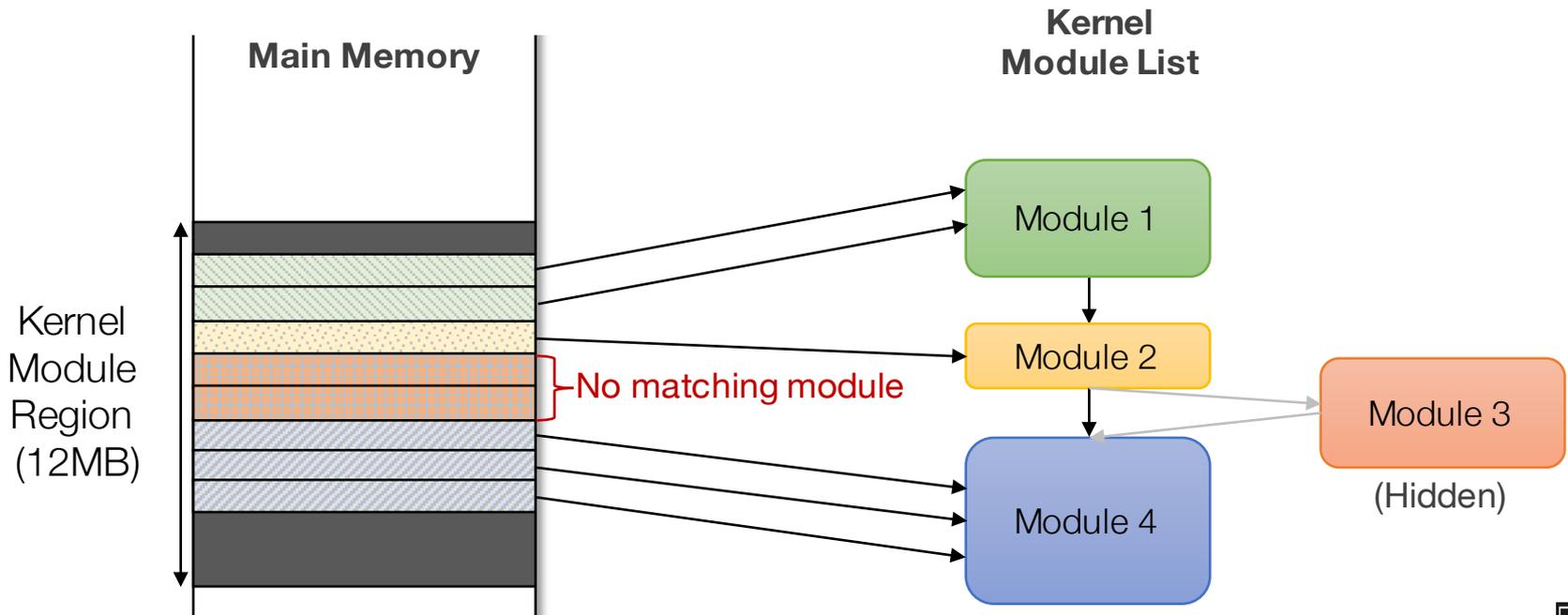
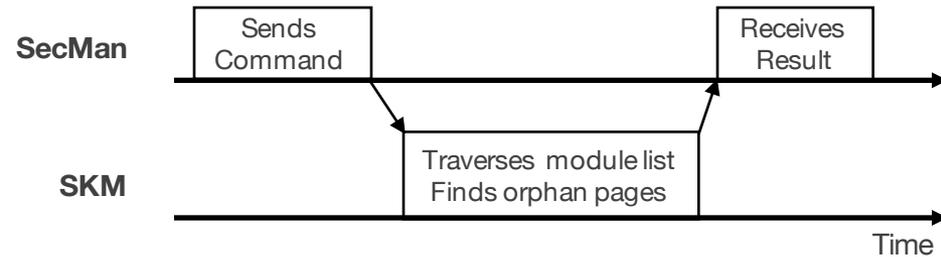
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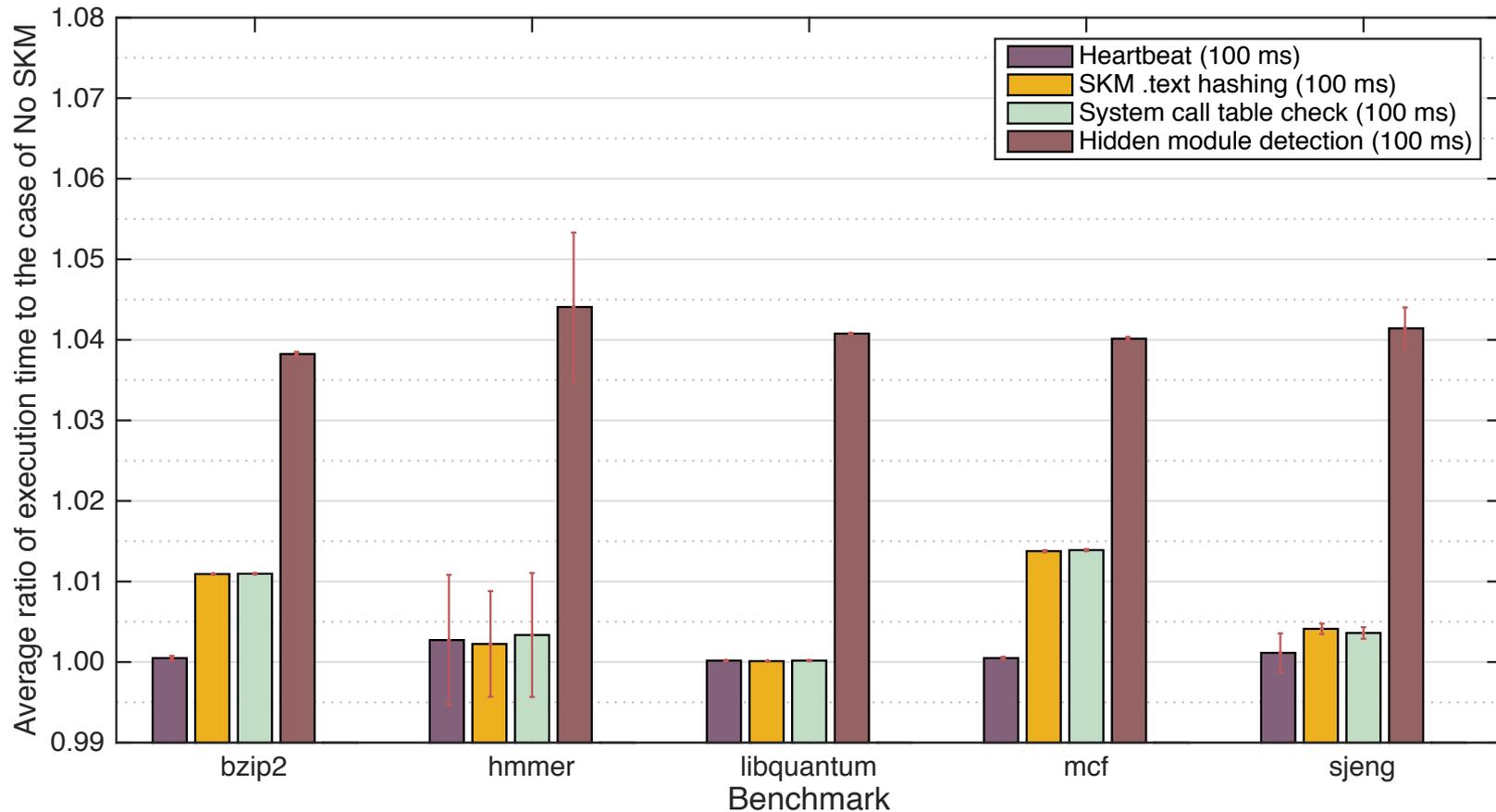
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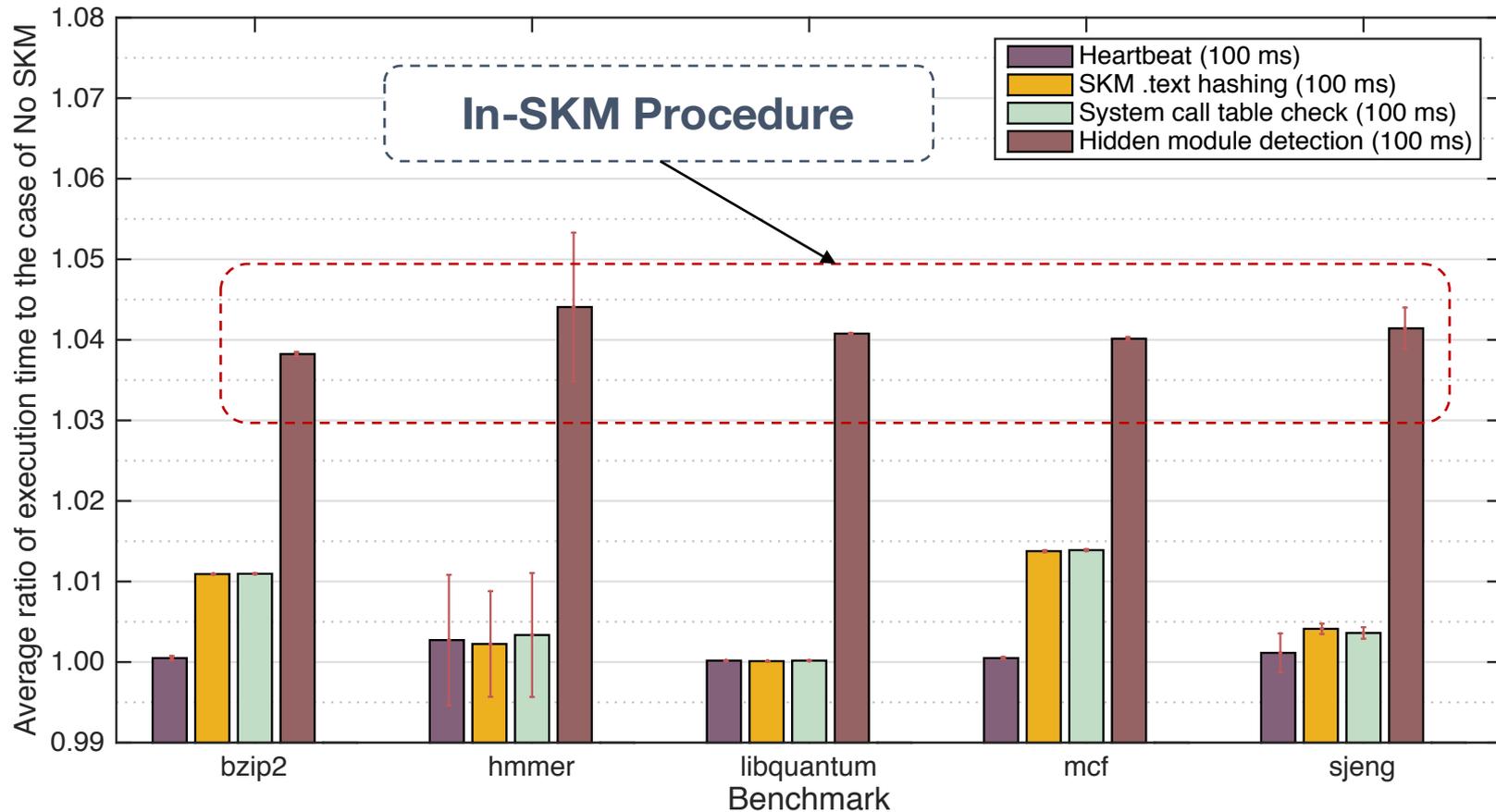
Evaluation – Performance Overhead

- SPEC Benchmarks on the monitored core



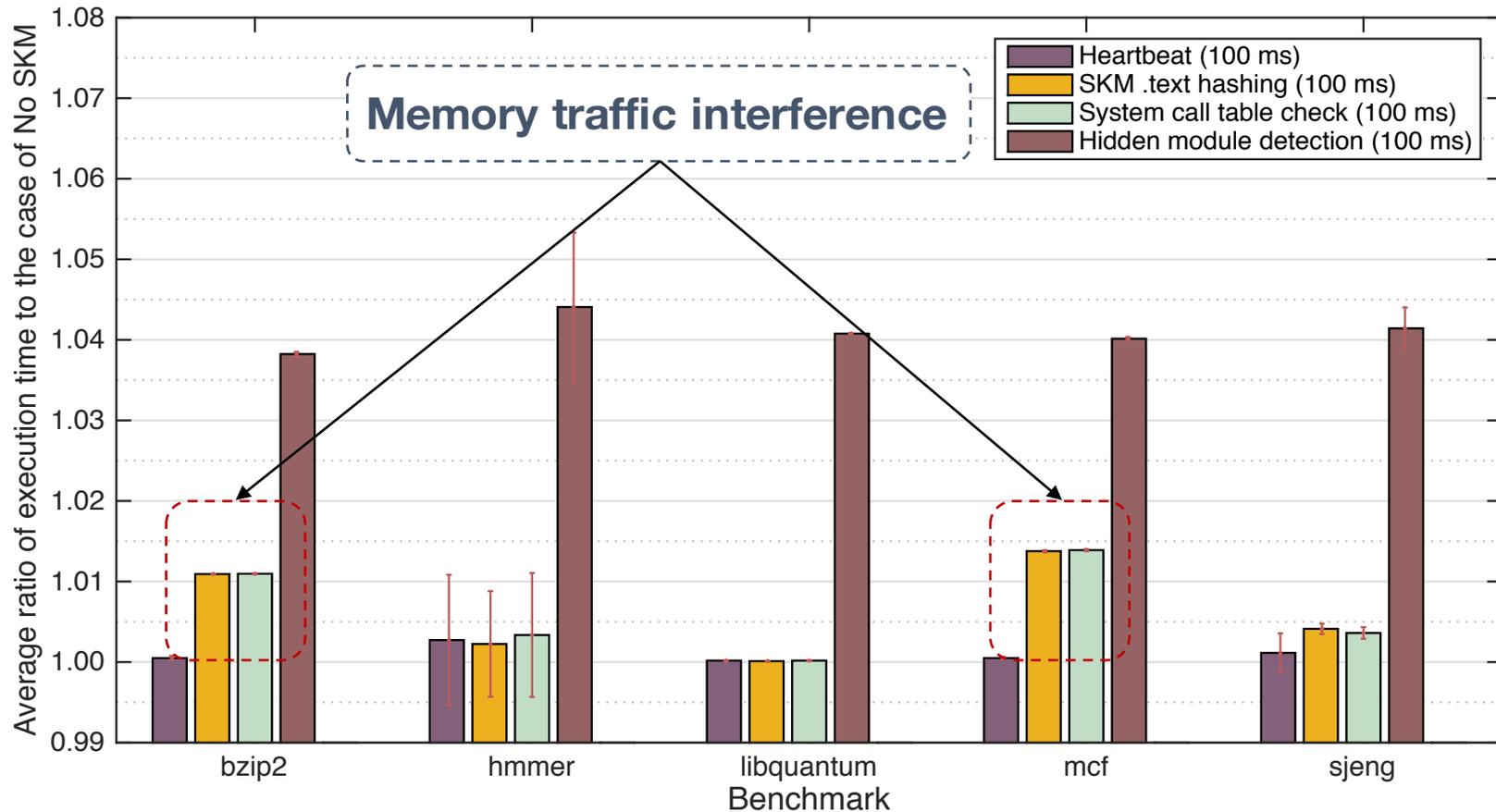
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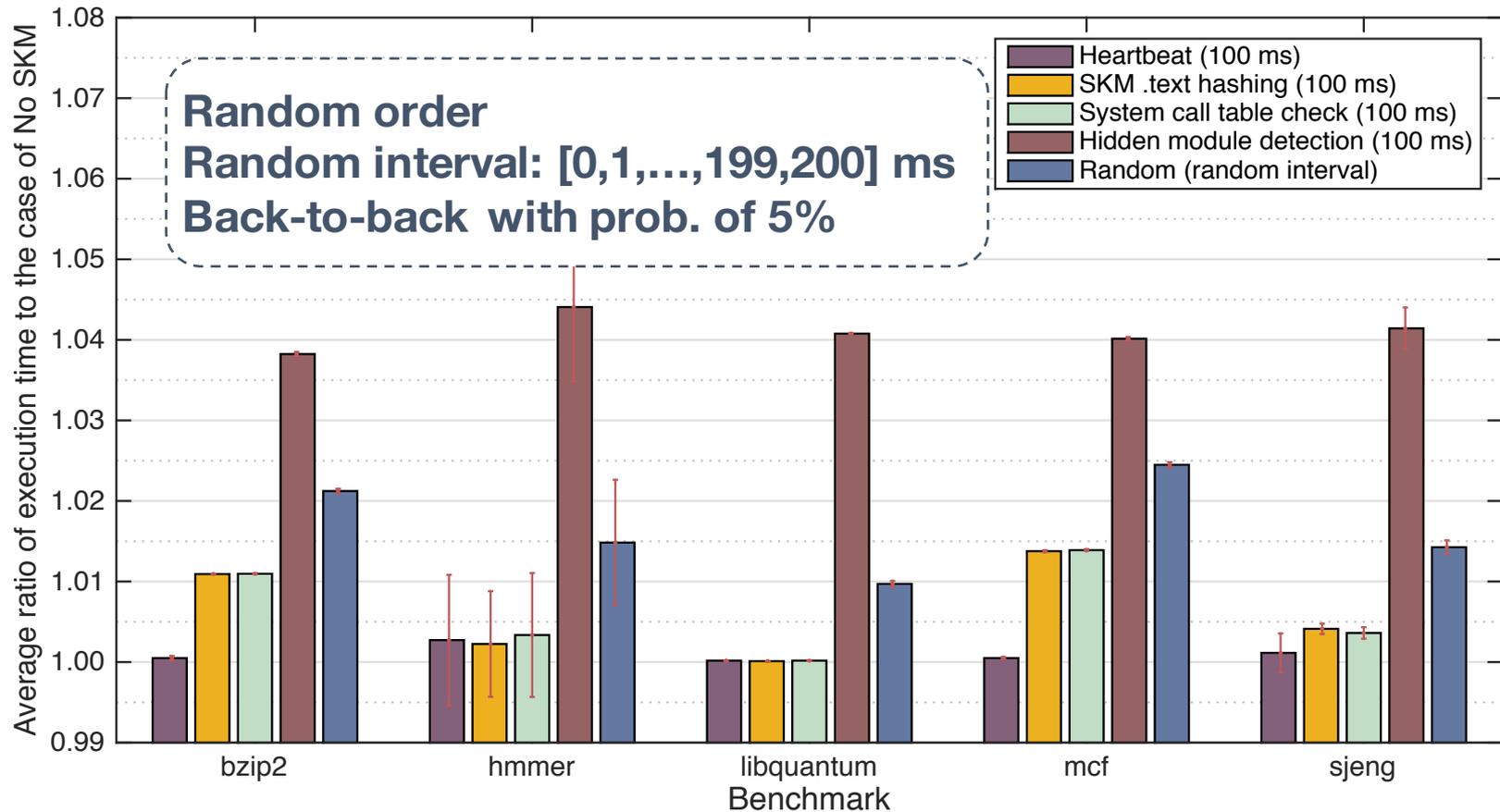
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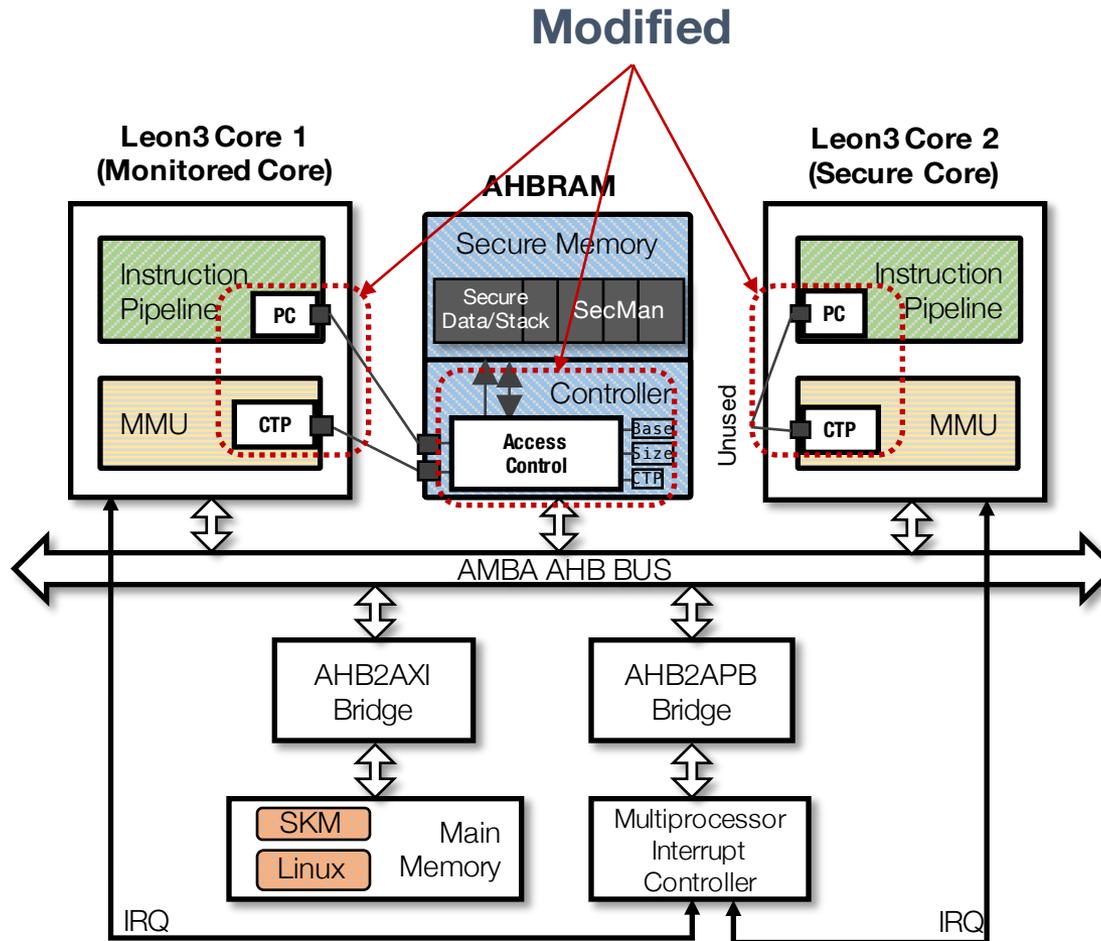


Evaluation – Performance Overhead

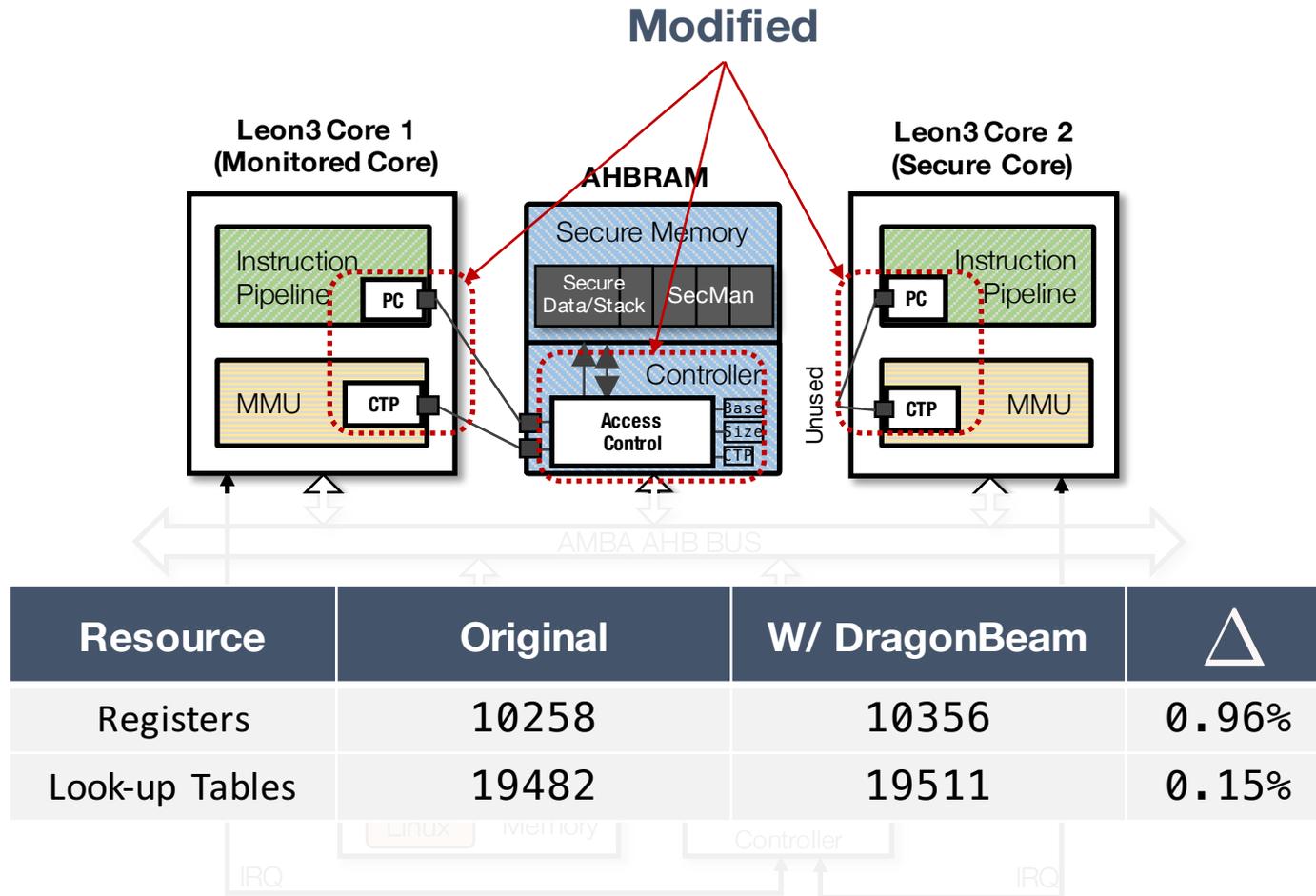
- SPEC Benchmarks on the monitored core



Evaluation – Hardware Cost

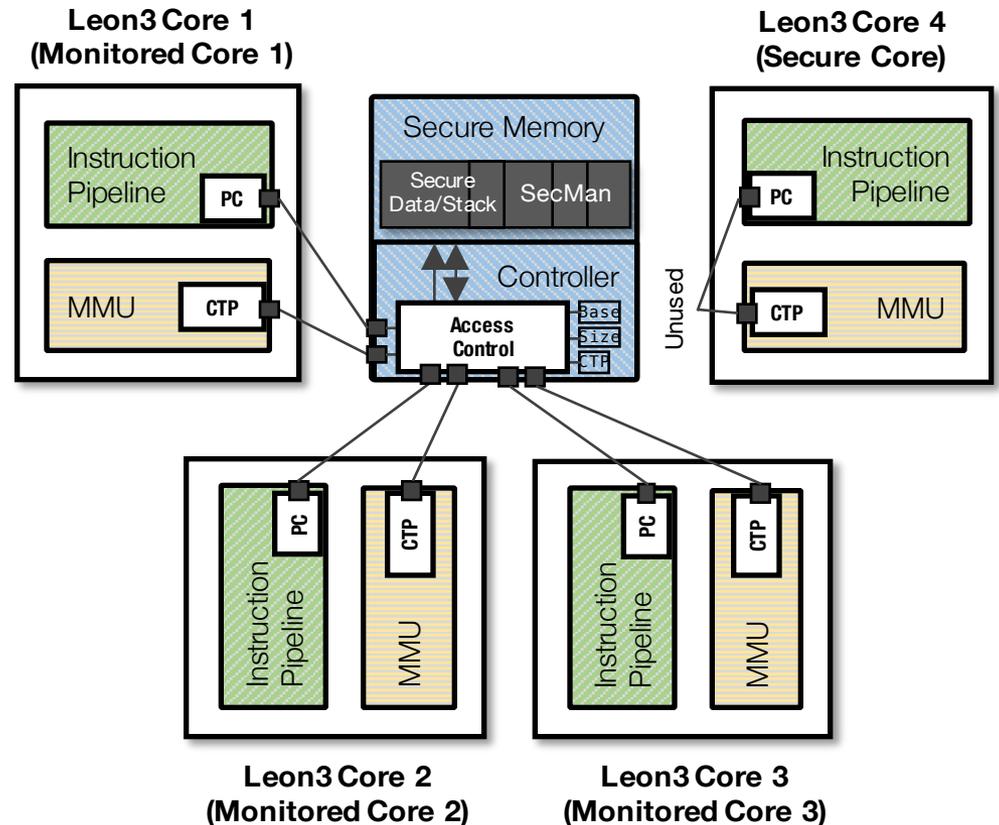


Evaluation – Hardware Cost



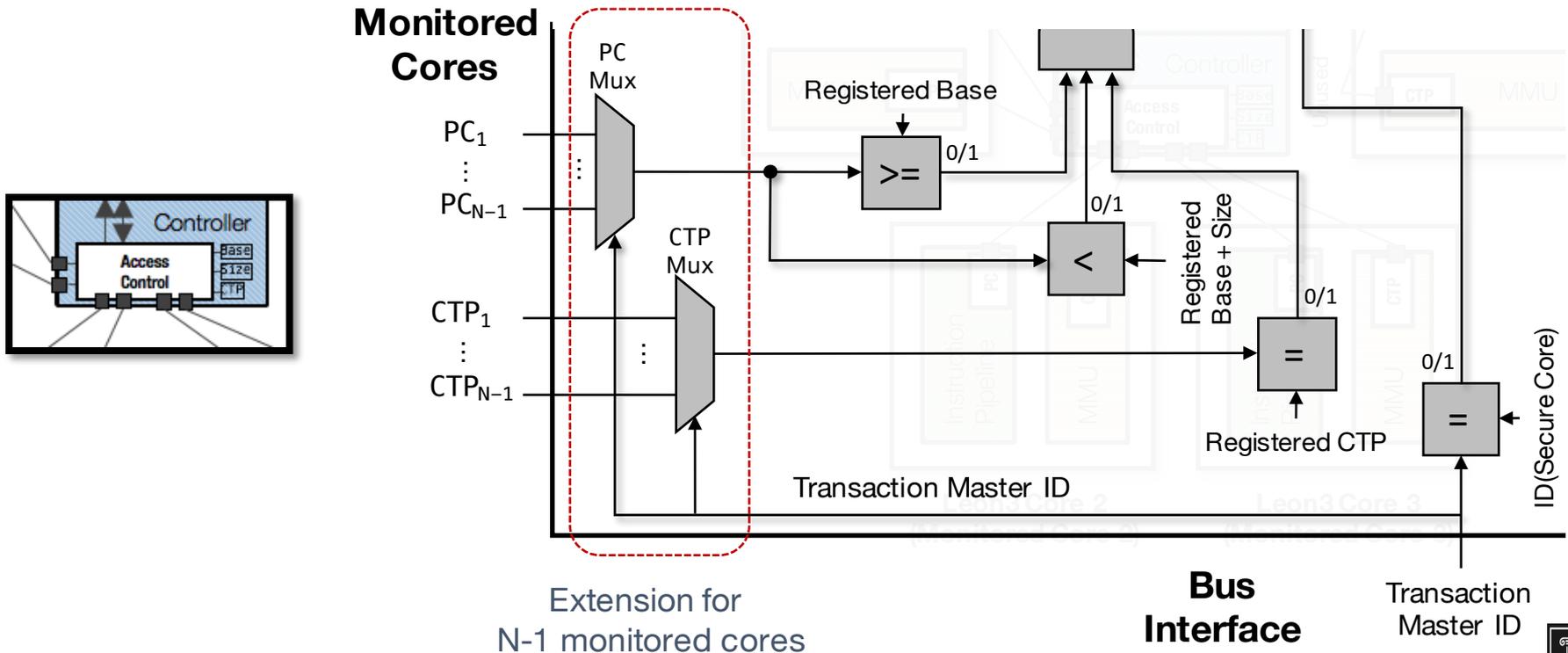
Extension to Multiple Monitored Cores

- Extended to quad-core
- Works only for SMP
 - Single SKM



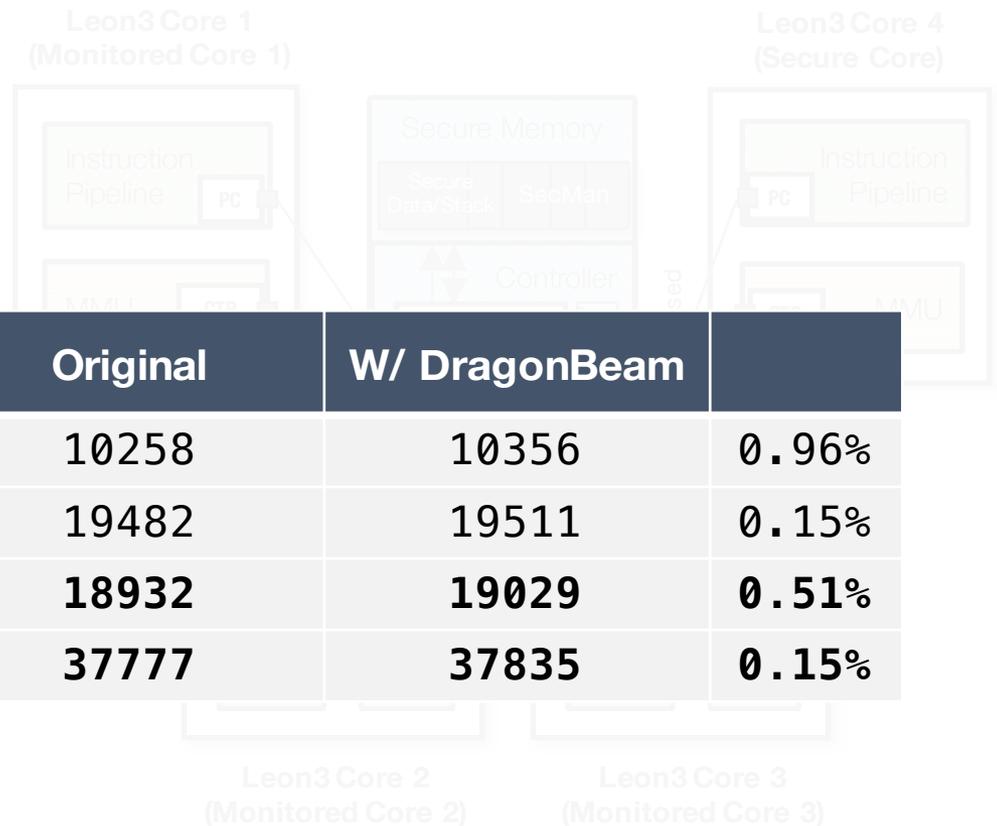
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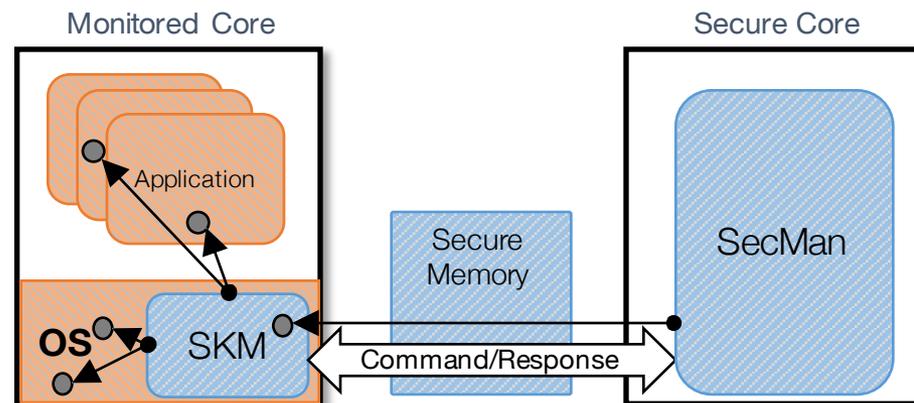
	Resource	Original	W/ DragonBeam	
Dual Core	Registers	10258	10356	0.96%
	Look-up Tables	19482	19511	0.15%
Quad Core	Registers	18932	19029	0.51%
	Look-up Tables	37777	37835	0.15%



Conclusion

- **DragonBeam Framework**

- HW/SW framework for trusted security monitoring
 - Bootstrap trust into SW layer from trusted HW
 - Multicore-based
- Expanded observability due to in-place monitoring
 - Secure “Kernel Module”
- Allows for customized security modules to system developers



Thank you

